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H8/300 Programming Manual

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Preface

The H8/300 CPU forms the common core of all chips in the H8/300 Series. Featuring a Hitachi-original, high-speed, RISC-like architecture, it has eight 16-bit (or sixteen 8-bit) general registers and a concise, optimized instruction set.

This manual gives detailed descriptions of the H8/300 instructions. The descriptions apply to all chips in the H8/300 Series. Assembly-language programmers should also read the separate *H8/300 Series Cross Assembler User's Manual*.

For hardware details, refer to the hardware manual of the specific chip.

Contents

Sec	tion 1. CPU	. 1
1.1	General CPU Architecture	. 2
1.2	Registers	. 5
1.3	Instructions	. 8
Sec	tion 2. Instruction Set	32
BCC	ADD (ADD binary) (byte)	
	ADD (ADD binary) (byte)	
	ADDS (ADD with Sign extension)	
	ADDX (ADD with eXtend carry)	
	AND (AND logical)	
	ANDC (AND Control register)	
	BAND (Bit AND)	
	Bcc (Branch conditionally)	
	•	
	BCLR (Bit CLeaR)	
	BIAND (Bit Invert Loop)	
	BILD (Bit Invert LoaD)	
	BIOR (Bit Invert OR)	
	BIST (Bit Invert STore)	
	BIXOR (Bit Invert eXclusive OR)	
	BLD (Bit LoaD)	
	BNOT (Bit NOT)	
	BOR (Bit inclusive OR)	
	BSET (Bit SET)	
	BSR (Branch to SubRoutine)	
	BST (Bit STore)	
	BTST (Bit TeST)	
	BXOR (Bit eXclusive OR)	
	CMP (CoMPare) (byte)	
	CMP (CoMPare) (word)	
	DAA (Decimal Adjust Add)	
	DAS (Decimal Adjust Subtract)	
	DEC (DECrement)	
	DIVXU (DIVide eXtend as Unsigned)	
	EEPMOV (MOVe data to EEPROM)	. 76

INC (INCrement)	78
JMP (JuMP)	79
JSR (Jump to SubRoutine)	80
LDC (LoaD to Control register)	81
MOV(MOVe data) (byte)	82
MOV(MOVe data) (word)	83
MOV(MOVe data) (byte)	84
MOV(MOVe data) (word)	85
MOV(MOVe data) (byte)	86
MOV(MOVe data) (word)	87
MOVFPE (MOVe data From Peripheral with E clock)	88
MOVTPE (MOVe data To Peripheral with E clock)	89
MULXU (MULtiply eXtend as Unsigned)	90
NEG (NEGate)	91
NOP (No OPeration)	92
NOT (NOT = logical complement)	93
OR (inclusive OR logical)	94
ORC (inclusive OR Control register)	95
POP (POP data)	96
PUSH (PUSH data)	97
ROTL (ROTate Left)	98
ROTR (ROTate Right)	99
ROTXL (ROTate with eXtend carry Left)	100
ROTXR (ROTate with eXtend carry Right)	101
RTE (ReTurn from Exception)	102
RTS (ReTurn from Subroutine)	103
SHAL (SHift Arithmetic Left)	104
SHAR (SHift Arithmetic Right)	105
SHLL (SHift Logical Left)	106
SHLR (SHift Logical Right)	107
SLEEP (SLEEP)	108
STC (STore from Control register)	109
SUB (SUBtract binary) (byte)	110
SUB (SUBtract binary) (word)	112
SUBS (SUBtract with Sign extension)	113
SUBX (SUBtract with eXtend carry)	114
XOR (eXclusive OR logical)	115

XORC (eX	Xclusive OR Control register)	116
Appendix A.	Operation Code Map	117
Appendix B.	Instruction Set List	118
Appendix C.	Number of Execution States	124

Section 1. CPU

This document is a reference manual for programming the H8/300, a high-speed central processing unit with a Hitachi-original RISC-like architecture that is employed as a CPU core in a series of low-cost single-chip microcomputers intended for applications ranging from smart cards to office and factory automation.

The H8/300 features a concise instruction set in which most frequently-used instructions are two bytes long and execute in just two states (0.2µs with a 10MHz system clock). Its general registers can be accessed as 16-bit word registers or 8-bit byte registers. The instruction set includes both 8-bit and 16-bit instructions.

Section 1 of this manual summarizes the CPU architecture and instruction set. Section 2 gives detailed descriptions of the instructions. Appendices give an operation code map, a complete list of the instruction set, and tables for calculating instruction execution time. Programmers should also refer to the *User's Manual* of the chip being programmed for information on bus cycles, interrupt service, I/O ports, power-down modes, and on-chip facilities such as memory and timers, and for a memory map.

1.1 General CPU Architecture

1.1.1 Features

Table 1-1 summarizes the CPU architecture. Figures 1-1 and 1-2 show how data are stored in registers and memory.

Table 1-1. CPU Architecture

Item	Description					
Address space	64K bytes, H'0000 to	64K bytes, H'0000 to H'FFFF				
Data types	Bit, 4-bit (packed Bo	CD), byte, word (2 bytes)				
General registers	Sixteen 8-bit general	l registers (R0H, R0L,, R7H, R7L),				
	ght 16-bit general registers (R0 to R7)					
Control registers	Program counter (PC	C)				
	Condition code regis	ster (CCR)				
Addressing modes	Rn	Register direct				
	@Rn	Register indirect				
	@(d:16, Rn)	Register indirect with 16-bit displacement				
	@Rn+	Register indirect with post-increment				
	@-Rn	Register indirect with pre-decrement				
	@aa:8, @aa:16	Absolute address (8 or 16 bits)				
	#xx:8, #xx:16	Immediate (8-, or 16-bit data)				
	@(d:8, PC)	PC-relative (8-bit displacement)				
	@@aa:8	Memory indirect				
Instruction length	struction length 2 or 4 bytes					

Notes:

- 1. Word data stored in memory must be stored at an even address.
- 2. Instructions must be stored at even addresses.
- 3. General register R7 is used as the stack pointer (SP).

1.1.2 Data Structure

The H8/300 CPU can process 1-bit data, 4-bit (packed BCD) data, 8-bit (byte) data, and 16-bit (word) data.

- Bit manipulation instructions operate on 1-bit data specified as bit n (n = 0, 1, 2, ..., 7) in a byte operand.
- All operational instructions except ADDS and SUBS can operate on byte data.

- The DAA and DAS instruction perform decimal arithmetic adjustments on byte data in packed BCD form. Each 4-bit of the byte is treated as a decimal digit.
- The MOV.W, ADD.W, SUB.W, CMP.W, ADDS, SUBS, MULXU (8 bits \times 8 bits), and DIVXU (16 bits \div 8 bits) instructions operate on word data.

Data Structure in General Registers: Data of all the sizes above can be stored in general registers as shown in figure 1-1.

Data type	Register No.	Data format
1-Bit data	RnH	7 0 7 6 5 4 3 2 1 0 Don't-care
1-Bit data	RnL	7 0 Don't-care 76543210
Byte data	RnH	7 0
Byte data	RnL	Don't-care
Word data	Rn	15 0
4-Bit BCD data	RnH	7 43 0 Upper digit Lower digit Don't-care
4-Bit BCD data	RnL	7 43 0 Don't-care Upper digit Lower digit

RnH: Upper 8 bits of General Register RnL: Lower 8 bits of General Register

MSB: Most Significant Bit LSB: Least Significant Bit

Figure 1-1. Register Data Structure

Memory Data Structure: Figure 1-2 indicates the data structure in memory.

Word data stored in memory must always begin at an even address. In word access the least significant bit of the address is regarded as "0." If an odd address is specified, no address error occurs but the access is performed at the preceding even address. This rule affects MOV.W instructions and branching instructions, and implies that only even addresses should be stored in the vector table.

Data type	Address	Data format
		7 0
1-Bit data	Address n	76543210
Byte data	Address n	M L S S B B B B B B B B B B B B B B B B B
Word data	Even address Odd address	Upper 8 bits Lower 8 bits
Byte data (CCR) on stack	Even address Odd address	S CCR S B L S CCR S B B CCR S B B CCR S B B B CCR S B B B CCR S B B CCR S B B CCR S B B B CCR S B B B CCR S B B CCR S B B CCR S B B B CCR S B B CCR S B B CCR S B B B CCR S B B CCR S B B CCR S B CCR S B B CCR S B CCR S B B CCR S B B CCR S B CCR S B B CCR S B B CCR S B B CCR S B CCR S B B CCR S B B CCR S B CCR S B CCR S B B CCR S B B CCR S B CCR S B B CCR S B B CCR S B CC
Word data on stack	Even address Odd address	S Upper 8 bits Lower 8 bits S

CCR: Condition code register.

Note: Word data must begin at an even address.

*: Ignored when return.

Figure 1-2. Memory Data Formats

The stack is always accessed a word at a time. When the CCR is pushed on the stack, two identical copies of the CCR are pushed to make a complete word. When they are returned, the lower byte is ignored.

1.1.3 Address Space

The H8/300 CPU supports a 64K-byte address space. The memory map differs depending on the particular chip in the H8/300 Series and its operating mode. See the *Hardware Manual* of the chip for details.

1.2 Registers

Figure 1-3 shows the register structure of the H8/300 CPU. There are sixteen 8-bit general registers (R0H, R0L, ..., R7H, R7L), which can also be accessed as eight 16-bit registers (R0 to R7). There are two control registers: the 16-bit program counter (PC) and the 8-bit condition code register (CCR).

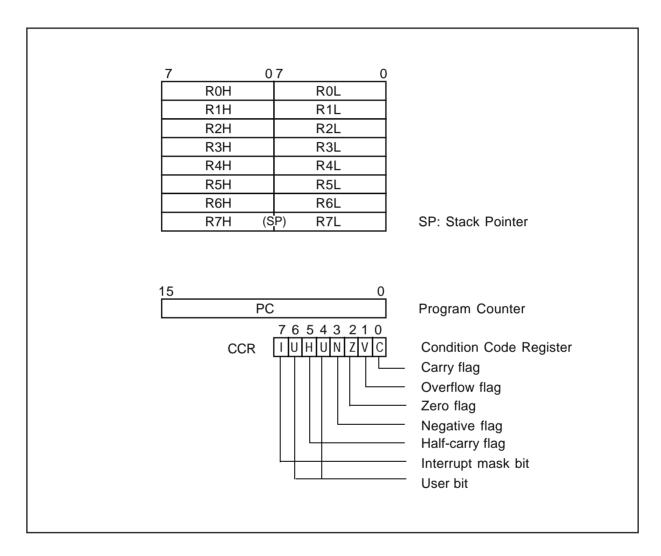


Figure 1-3. CPU Registers

1.2.1 General Registers

All the general registers can be used as both data registers and address registers. When used as address registers, the general registers are accessed as 16-bit registers (R0 to R7). When used as data registers, they can be accessed as 16-bit registers (R0 to R7), or the high (R0H to R7H) and low (R0L to R7L) bytes can be accessed separately as 8-bit registers. The register length is determined by the instruction.

R7 also functions as the stack pointer, used implicitly by hardware in processing interrupts and subroutine calls. In assembly language, the letters SP can be coded as a synonym for R7. As indicated in figure 1-4, R7 (SP) points to the top of the stack.

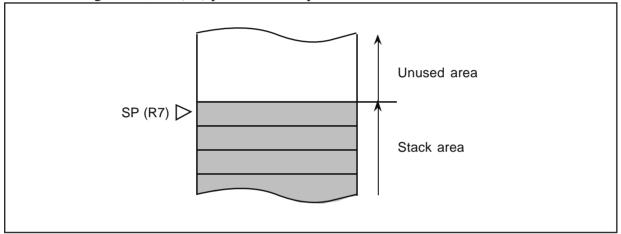


Figure 1-4. Stack Pointer

1.2.2 Control Registers

The CPU has a 16-bit program counter (PC) and an 8-bit condition code register (CCR).

- (1) **Program Counter (PC):** This 16-bit register indicates the address of the next instruction the CPU will execute. Instructions are fetched by 16-bit (word) access, so the least significant bit of the PC is ignored (always regarded as 0).
- (2) Condition Code Register (CCR): This 8-bit register indicates the internal status of the CPU with an interrupt mask (I) bit and five flag bits: half-carry (H), negative (N), zero (Z), overflow (V), and carry (C) flags. The two unused bits are available to the user. The bit configuration of the condition code register is shown below.

Bit	7	6	5	4	3	2	1	0
	Ι	U	Н	U	N	Z	V	C
Initial value	1	*	*	*	*	*	*	*
Read/Write	R/W							

 ^{*} Undetermined

Bit 7—Interrupt Mask Bit (I): When this bit is set to "1," all interrupts except NMI are masked. This bit is set to "1" automatically by a reset and at the start of interrupt handling.

Bits 6 and 4—User Bits (U): These bits can be written and read by software for its own purposes.

Bit 5—Half-Carry (H): This bit is used by add, subtract, and compare instructions to indicate a borrow or carry out of bit 3 or bit 11. It is referenced by the decimal adjust instructions.

Bit 3—Negative (N): This bit indicates the most significant bit (sign bit) of the result of an instruction.

Bit 2—Zero (Z): This bit is set to "1" to indicate a zero result and cleared to "0" to indicate a nonzero result.

Bit 1—Overflow (V): This bit is set to "1" when an arithmetic overflow occurs, and cleared to "0" at other times.

Bit 0—Carry (C): This bit is used by:

- Add, subtract, and compare instructions, to indicate a carry or borrow at the most significant bit
- Shift and rotate instructions, to store the value shifted out of the most or least significant bit
- Bit manipulation instructions, as a bit accumulator

System control instructions can load and store the CCR, and perform logic operations to set, clear, or toggle selected bits.

1.2.3 Initial Register Values

When the CPU is reset, the program counter (PC) is loaded from the vector table and the interrupt mask bit (I) in the CCR is set to "1." The other CCR bits and the general registers are not initialized.

In particular, the stack pointer (R7) is not initialized. To prevent program crashes the stack pointer should be initialized by software, by the first instruction executed after a reset.

1.3 Instructions

Features:

- The H8/300 has a concise set of 57 RISC-like instructions.
- Arithmetic and logic are performed as register-to-register operations, or with immediate data.
- All instructions are 2 or 4 bytes long.
- Fast multiply/divide instructions; extensive bit manipulation instructions.
- Eight addressing modes.

1.3.1 Types of Instructions

Table 1-2 classifies the H8/300 instructions by type. Tables 1-3 to 1-10 briefly describe their functions. Section 2, Instruction Set, gives detailed descriptions.

Table 1-2. Instruction Classification

Function	Instructions	ypes
Data transfer	MOV, MOVFPE, MOVTPE, POP*, PUSH*	3
Arithmetic operation	$_{ m S}$ ADD, SUB, ADDX, SUBX, INC, DEC, ADDS, SUBS,	14
	DAA, DAS, MULXU, DIVXU, CMP, NEG	
Logic operations	AND, OR, XOR, NOT	4
Shift	SHAL, SHAR, SHLL, SHLR, ROTL, ROTR, ROTXL,	8
	ROTXR	
Bit manipulation	BSET, BCLR, BNOT, BTST, BAND, BIAND, BOR	14
	BIOR, BXOR, BIXOR, BLD, BILD, BST, BIST	
Branch	Bcc**, JMP, BSR, JSR, RTS	5
System control	RTE, SLEEP, LDC, STC, ANDC, ORC, XORC, NOP	8
Block data transfer	EEPMOV	1
	Total	57

^{*} POP Rn is equivalent to MOV.W @SP+, Rn. PUSH Rn is equivalent to MOV.W Rn, @-SP.

^{**} Bcc is a conditional branch instruction in which cc represents a condition .

1.3.2 Instruction Functions

Notation

Tables 1-3 to 1-10 give brief descriptions of the instructions in each functional group. The following notation is used.

Rd	General register (destination)		
Rs	General register (source)		
Rn	General register		
(EAd)	Destination operand		
(EAs)	Source operand		
CCR	Condition code register		
N	N (negative) bit of CCR		
Z	Z (zero) bit of CCR		
V	V (overflow) bit of CCR		
С	C (carry) bit of CCR		
PC	Program counter		
SP	Stack pointer (R7)		
#Imm	Immediate data		
#xx:3	3-Bit immediate data		
#xx:8	8-Bit immediate data		

#xx:16 16-Bit immediate data
op Operation field

Displacement

Addition

Division
AND logical
OR logical

Move

Not

Subtraction Multiplication

op disp

<u>+</u> __

× ÷ ^ ∨ ⊕

:3, :8, :16 3-bit, 8-bit, or 16-bit length.

Exclusive OR logical

Table 1-3. Data Transfer Instructions

Instruction	Size*	Function
MOV	B/W	$(EAs) \rightarrow Rd, Rs \rightarrow (EAd)$
		Moves data between two general registers or between a general
		register and memory, or moves immediate data to a general register.
		The Rn, @Rn, @(d:16, Rn), @aa:16, #xx:8 or #xx:16, @-Rn, and
		@Rn+ addressing modes are available for byte or word data. The
		@aa:8 addressing mode is available for byte data only.
		The @-R7 and @R7+ modes require word operands. Do not
		specify byte size for these two modes.
MOVFPE	В	(EAs) Rd
		Transfers data from memory to a general register in
		synchronization with the E clock.
MOVTPE	В	Rs \rightarrow (EAd)
		Transfers data from a general register to memory in
		synchronization with the E clock.
POP	W	$@SP+ \rightarrow Rn$
		Pops a 16-bit general register from the stack.
		Equivalent to MOV.W @SP+, Rn.
PUSH	W	$Rn \rightarrow @-SP$
		Pushes a 16-bit general register onto the stack.
		Equivalent to MOV.W Rn, @-SP.

B: Byte W: Word

Table 1-4. Arithmetic Instructions

Instruction	Size*	Function
ADD	B/W	$Rd \pm Rs \rightarrow Rd, Rd + \#Imm \rightarrow Rd$
SUB		Performs addition or subtraction on data in two general registers,
		or addition on immediate data and data in a general register.
		Immediate data cannot be subtracted from data in a general register.
		Word data can be added or subtracted only when both words are in
		general registers.
ADDX	В	$Rd \pm Rs \pm C \rightarrow Rd, Rd \pm \#Imm \pm C \rightarrow Rd$
SUBX		Performs addition or subtraction with carry or borrow on byte data
		in two general registers, or addition or subtraction on immediate data
		and data in a general register.
INC	В	$Rd \pm 1 \rightarrow Rd$
DEC		Increments or decrements a general register.
ADDS	W	$Rd \pm 1 \rightarrow Rd, Rd \pm 2 \rightarrow Rd$
SUBS		Adds or subtracts immediate data to or from data in a general
		register. The immediate data must be 1 or 2.
DAA	В	Rd decimal adjust \rightarrow Rd
DAS		Decimal-adjusts (adjusts to packed BCD) an addition or
		subtraction result in a general register by referring to the CCR.
MULXU	В	$Rd \times Rs \rightarrow Rd$
		Performs 8-bit \times 8-bit unsigned multiplication on data in two
		general registers, providing a 16-bit result.
DIVXU	В	$Rd \div Rs \longrightarrow Rd$
		Performs 16-bit ÷ 8-bit unsigned division on data in two general
		registers, providing an 8-bit quotient and 8-bit remainder.
CMP	B/W	Rd - Rs, $Rd - #Imm$
		Compares data in a general register with data in another general
		register or with immediate data. Word data can be compared only
		between two general registers.
NEG	В	$0-Rd \rightarrow Rd$
		Obtains the two's complement (arithmetic complement) of data in
		a general register.

B: Byte W: Word

Table 1-5. Logic Operation Instructions

Instruction	Size*	Function
AND	В	$Rd \ _{\wedge} \ Rs \ _{\rightarrow} \ Rd, Rd \ _{\wedge} \ \#Imm \ _{\rightarrow} \ Rd$
		Performs a logical AND operation on a general register and
		another general register or immediate data.
OR	В	Rd $_{\vee}$ Rs $_{\rightarrow}$ Rd, Rd $_{\vee}$ #Imm $_{\rightarrow}$ Rd
		Performs a logical OR operation on a general register and another
		general register or immediate data.
XOR	В	$Rd \oplus Rs \rightarrow Rd$, $Rd \oplus \#Imm \rightarrow Rd$
		Performs a logical exclusive OR operation on a general register
		and another general register or immediate data.
NOT	В	$\neg \operatorname{Rd} \rightarrow \operatorname{Rd}$
		Obtains the one's complement (logical complement) of general
		register contents.

B: Byte

Table 1-6. Shift Instructions

Instruction	Size*	Function
SHAL	В	$Rd ext{ shift } \rightarrow Rd$
SHAR		Performs an arithmetic shift operation on general register contents.
SHLL	В	$Rd ext{ shift } \rightarrow Rd$
SHLR		Performs a logical shift operation on general register contents.
ROTL	В	Rd rotate \rightarrow Rd
ROTR		Rotates general register contents.
ROTXL	В	Rd rotate through carry \rightarrow Rd
ROTXR		Rotates general register contents through the C (carry) bit.

* Size: Operand size

B: Byte

Table 1-7. Bit-Manipulation Instructions

Instruction	Size*	Function
BSET	В	$1 \rightarrow (\langle \text{bit-No.} \rangle \text{ of } \langle \text{EAd} \rangle)$
		Sets a specified bit in a general register or memory to "1." The bit
		is specified by a bit number, given in 3-bit immediate data or the lower
		three bits of a general register.
BCLR	В	$0 \rightarrow (\langle \text{bit-No.} \rangle \text{ of } \langle \text{EAd} \rangle)$
		Clears a specified bit in a general register or memory to "0." The
		bit is specified by a bit number, given in 3-bit immediate data or the
		lower three bits of a general register.
BNOT	В	$\neg(\langle \text{bit-No.}\rangle \text{ of } \langle \text{EAd}\rangle) \rightarrow (\langle \text{bit-No.}\rangle \text{ of } \langle \text{EAd}\rangle)$
		Inverts a specified bit in a general register or memory. The bit is
		specified by a bit number, given in 3-bit immediate data or the lower
		three bits of a general register.
BTST	В	$\neg (< bit-No. > of < EAd >) \rightarrow Z$
		Tests a specified bit in a general register or memory and sets or
		clears the Z flag accordingly. The bit is specified by a bit number,
		given in 3-bit immediate data or the lower three bits of a general
		register.
BAND	В	$C \land (< bit-No. > of < EAd >) \rightarrow C$
		ANDs the C flag with a specified bit in a general register or
		memory.
BIAND	В	$C \cap [\neg (< bit-No. > of < EAd >)] \rightarrow C$
		ANDs the C flag with the inverse of a specified bit in a general
		register or memory.
		The bit number is specified by 3-bit immediate data.
BOR	В	$C \downarrow (< bit-No. > of < EAd >) \rightarrow C$
		ORs the C flag with a specified bit in a general register or memory.
BIOR	В	$C \setminus [\neg (\langle bit\text{-No.} \rangle \text{ of } \langle EAd \rangle)] \rightarrow C$
		ORs the C flag with the inverse of a specified bit in a general
		register or memory.
		The bit number is specified by 3-bit immediate data.

Table 1-7. Bit-Manipulation Instructions (Cont.)

Instruction	Size*	Function
BXOR	В	$C \oplus (\langle \text{bit-No.} \rangle \text{ of } \langle \text{EAd} \rangle) \rightarrow C$
		Exclusive-ORs the C flag with a specified bit in a general register
		or memory.
BIXOR	В	$C \oplus [\neg (of)] \rightarrow C$
		Exclusive-ORs the C flag with the inverse of a specified bit in a
		general register or memory.
		The bit number is specified by 3-bit immediate data.
BLD	В	$(<$ bit-No. $>$ of $<$ EAd $>) \rightarrow C$
		Copies a specified bit in a general register or memory to the C flag.
BILD	В	$\neg (< bit-No. > of < EAd >) \rightarrow C$
		Copies the inverse of a specified bit in a general register or
		memory to the C flag.
		The bit number is specified by 3-bit immediate data.
BST	В	$C \rightarrow (< bit-No. > of < EAd >)$
		Copies the C flag to a specified bit in a general register or memory.
BIST	В	$\neg C \rightarrow (< bit-No. > of < EAd >)$
		Copies the inverse of the C flag to a specified bit in a general
		register or memory.
		The bit number is specified by 3-bit immediate data.

B: Byte

Table 1-8. Branching Instructions

Instruction	Size	Function								
Всс	_	Branches if co	Branches if condition cc is true.							
		Mnemonic	cc Field	Description	Condition					
		BRA (BT)	0000	Always (True)	Always					
		BRN (BF)	0001	Never (False)	Never					
		BHI	0010	High	$C_{\vee}Z=0$					
		BLS	0011	Low or Same	$C_{\vee}Z=1$					
		BCC (BHS)	0100	Carry Clear	C = 0					
				(High or Same)						
		BCS (BLO)	0101	Carry Set (Low)	C = 1					
		BNE	0110	Not Equal	Z = 0					
		BEQ	0111	Equal	Z = 1					
		BVC	$1\ 0\ 0\ 0$	Overflow Clear	V = 0					
		BVS	1001	Overflow Set	V = 1					
		BPL	1010	Plus	N = 0					
		BMI	1011	Minus	N = 1					
		BGE	1100	Greater or Equal	$N \oplus V = 0$					
		BLT	1 1 0 1	Less Than	$N \oplus V = 1$					
		BGT	1110	Greater Than	$Z_{\vee}(N_{\bigoplus}V)=0$					
		BLE	1111	Less or Equal	$Z_{\vee}(N_{\bigoplus}V)=1$					
JMP		Branches unce	onditionally to	o a specified address	SS.					
BSR		Branches to a	subroutine at	a specified address	S.					
JSR		Branches to a	subroutine at	a specified displac	ement from the current					
		address.								
RTS	_	Returns from	a subroutine.							

Table 1-9. System Control Instructions

Instruction	Size*	Function					
RTE	_	Returns from an exception-handling routine.					
SLEEP	_	Causes a transition to the power-down state.					
LDC	В	$Rs \rightarrow CCR, \#Imm \rightarrow CCR$					
		Moves immediate data or general register contents to the condition					
		code register.					
STC	В	$CCR \rightarrow Rd$					
		Copies the condition code register to a specified general register.					
ANDC	В	$CCR \wedge \#Imm \rightarrow CCR$					
		Logically ANDs the condition code register with immediate data.					
ORC	В	$CCR \downarrow \#Imm \rightarrow CCR$					
		Logically ORs the condition code register with immediate data.					
XORC	В	$CCR \oplus \#Imm CCR$					
		Logically exclusive-ORs the condition code register with immediate					
		data.					
NOP		$PC + 2 \rightarrow PC$					
		Only increments the program counter.					

B: Byte

Table 1-10. Block Data Transfer Instruction

Instruction	Size	Function
EEPMOV	_	if $R4L \neq 0$ then
		repeat $@R5+ \rightarrow @R6+$
		R4L-1 R4L
		until $R4L = 0$
		else next;
		Moves a data block according to parameters set in general registers
		R4L, R5, and R6.
		R4L: size of block (bytes)
		R5: starting source address
		R6: starting destination address
		Execution of the next instruction starts as soon as the block transfer is
		completed.

Notes on Bit Manipulation Instructions: BSET, BCLR, BNOT, BST, and BIST are read-modify-write instructions. They read a byte of data, modify one bit in the byte, then write the byte back. Care is required when these instructions are applied to registers with write-only bits and to the I/O port registers.

Se	equence	Operation
1	Read	Read one data byte at the specified address
2	Modify	Modify one bit in the data byte
3	Write	Write the modified data byte back to the specified address

Example 1: BCLR is executed to clear bit 0 in the port 4 data direction register (P4DDR) under the following conditions.

P47: Input pin, Low, MOS pull-up transistor on

P46: Input pin, High, MOS pull-up transistor off

P45 – P40: Output pins, Low

The intended purpose of this BCLR instruction is to switch P40 from output to input.

Before Execution of BCLR Instruction

	P47	P46	P45	P44	P43	P42	P41	P40
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low	High	Low	Low	Low	Low	Low	Low
DDR	0	0	1	1	1	1	1	1
DR	1	0	0	0	0	0	0	0
Pull-up	On	Off	Off	Off	Off	Off	Off	Off

Execution of BCLR Instruction

BCLR #0 @P4DDR ; clear bit 0 in data direction register

After Execution of BCLR Instruction

	P47	P46	P45	P44	P43	P42	P41	P40
Input/output	Output	Input						
Pin state	Low	High	Low	Low	Low	Low	Low	High
DDR	1	1	1	1	1	1	1	0
DR	1	0	0	0	0	0	0	0
Pull-up	Off	Off						

Explanation: To execute the BCLR instruction, the CPU begins by reading P4DDR. Since P4DDR is a write-only register, it is read as H'FF, even though its true value is H'3F.

Next the CPU clears bit 0 of the read data, changing the value to H'FE.

Finally, the CPU writes this value (H'FE) back to P4DDR to complete the BCLR instruction.

As a result, P40DDR is cleared to "0," making P40 an input pin. In addition, P47DDR and P46DDR are set to "1," making P47 and P46 output pins.

Example 2: BSET is executed to set bit 0 in the port 4 data register (P4DR) under the following conditions.

P47: Input pin, Low, MOS pull-up transistor on

P46: Input pin, High, MOS pull-up transistor off

P45 – P40: Output pins, Low

The intended purpose of this BSET instruction is to switch the output level at P4⁰ from Low to High.

Before Execution of BSET Instruction

	P47	P46	P45	P44	P43	P42	P41	P40
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low	High	Low	Low	Low	Low	Low	Low
DDR	0	0	1	1	1	1	1	1
DR	1	0	0	0	0	0	0	0
Pull-up	On	Off	Off	Off	Off	Off	Off	Off

Execution of BSET Instruction

BSET #0 @PORT4 ; set bit 0 in port-4 data register

After Execution of BSET Instruction

	P47	P46	P45	P44	P43	P42	P41	P40
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low	High	Low	Low	Low	Low	Low	High
DDR	0	0	1	1	1	1	1	1
DR	0	1	0	0	0	0	0	1
Pull-up	Off	On	Off	Off	Off	Off	Off	Off

Explanation: To execute the BSET instruction, the CPU begins by reading port 4. Since P47 and P46 are input pins, the CPU reads the level of these pins directly, not the value in the data register. It reads P47 as Low ("0") and P46 as High ("1").

Since P4⁵ to P4⁰ are output pins, for these pins the CPU reads the value in the data register ("0"). The CPU therefore reads the value of port 4 as H'40, although the actual value in P4DR is H'80.

Next the CPU sets bit 0 of the read data to "1," changing the value to H'41.

Finally, the CPU writes this value (H'41) back to P4DR to complete the BSET instruction.

As a result, bit P40 is set to "1," switching pin P40 to High output. In addition, bits P47 and P46 are both modified, changing the on/off settings of the MOS pull-up transistors of pins P47 and P46.

Programming Solution: The switching of the pull-ups for P4⁷ and P4⁶ in example 2 can be avoided by storing the same data in both the port-4 data register and in a work area in RAM. Bit manipulations are performed on the data in the work area, after which the result is moved into the port-4 data register. In the following example RAM0 is a symbol for the user-selected address of the work area.

Before Execution of BSET Instruction

MOV.B	#80	R0L	; write data (H'80) for data register
MOV.B	R0L	@RAM0	; write to DR work area (RAM0)
MOV.B	R0L	@PORT4	; write to DR

	P4 ⁷	P46	P4 ⁵	P4 ⁴	P4 ³	P4 ²	P4 ¹	P40
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low	High	Low	Low	Low	Low	Low	Low
DDR	0	0	1	1	1	1	1	1
DR	1	0	0	0	0	0	0	0
Pull-up	On	Off	Off	Off	Off	Off	Off	Off
RAM0	1	0	0	0	0	0	0	0

Execution of BSET Instruction

BSET #0 @RAMO ; set bit 0 in DR work area (RAMO)

After Execution of BSET Instruction

MOV.B @RAMO ROL ; get value in work area (RAMO)

MOV.B ROL @PORT4 ; write value to DR

	P47	P46	P45	P44	P43	P42	P41	P40
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low	High	Low	Low	Low	Low	Low	High
DDR	0	0	1	1	1	1	1	1
DR	1	0	0	0	0	0	0	1
Pull-up	On	Off	Off	Off	Off	Off	Off	Off
RAM0	1	0	0	0	0	0	0	1

1.3.3 Machine-Language Coding

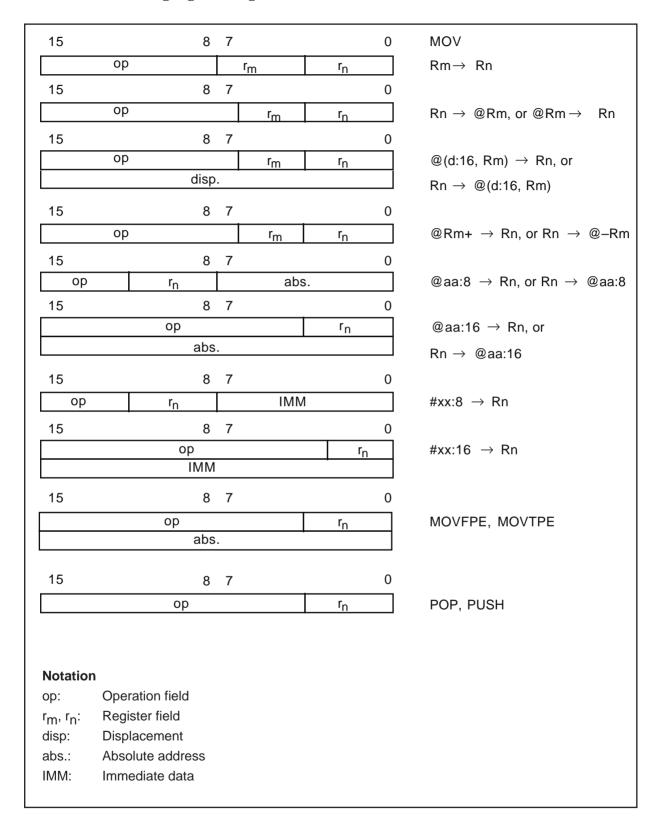


Figure 1-5. Machine-Language Coding of Data Transfer Instructions

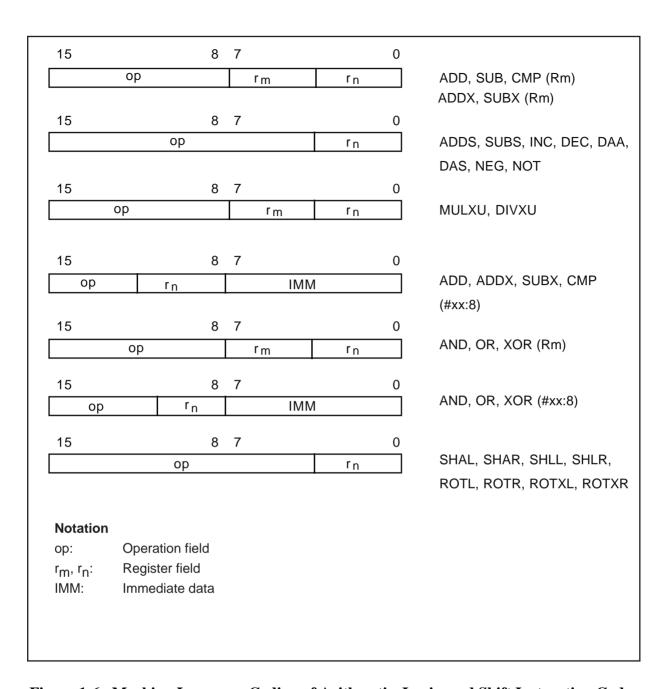


Figure 1-6. Machine-Language Coding of Arithmetic, Logic, and Shift Instruction Codes

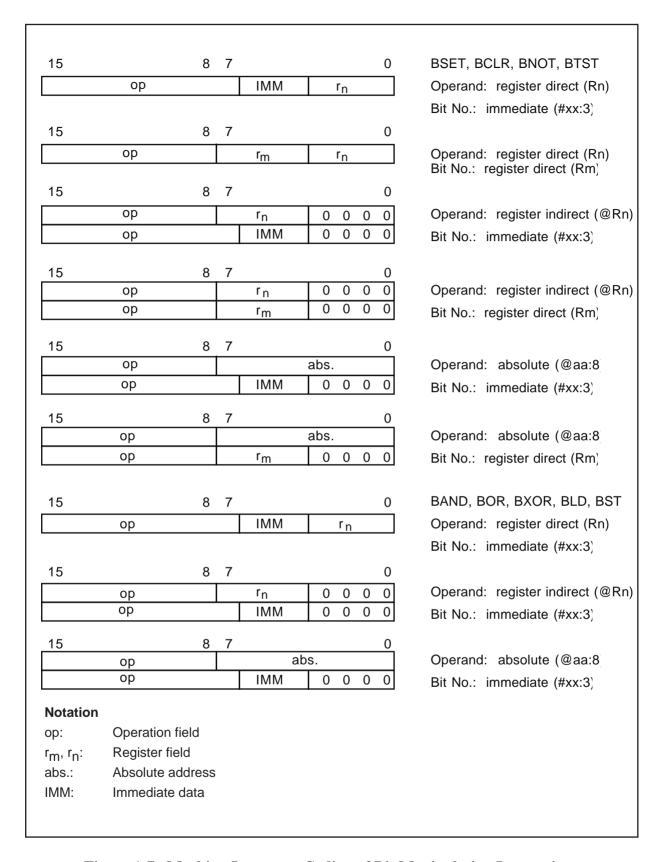


Figure 1-7. Machine-Language Coding of Bit Manipulation Instructions

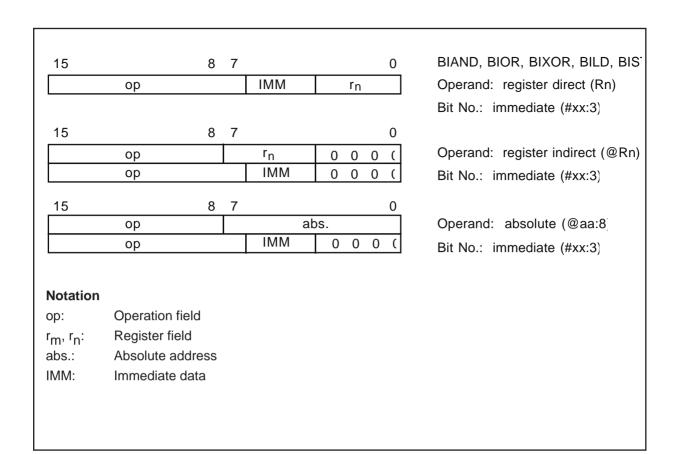


Figure 1-7. Machine-Language Coding of Bit Manipulation Instructions (Cont.)

15	:	8	7				0	
ор	cc			disp				Bcc
15	;	8	7				0	
	ор			r _m	0	0 0	(JMP (@Rm)
15		8	7				0	
		ор						JMP (@aa:16)
		abs).					,
15		8	7				0	
	ор			abs				JMP (@@aa:8
15	:	8	7				0	
	ор			dis	p.			BSR
15	;	8	7				0	
	ор			r _m	0	0 0	0	JSR (@Rm)
15		8	7				0	
		ор						JSR (@aa:16)
		abs						
15	:	8	7				0	
	ор			abs				JSR (@@aa:8
15	;	8	7				0	
ор						RTS		
N. 4.41								
Notatio								
op: cc:	Operation field Condition field							
	Register field							
r _m : disp.:	Displacement							
abs.:	Absolute addres	s						

and to the I/O port registers.

Example 1: BCLR is executed to clear bit 0 in the port 4 data direction register (P4DDR) under the following conditions.

P47: Input pin, Low, MOS pull-up transistor on

P46: Input pin, High, MOS pull-up transistor off

 $P4^5 - P4^0$: Output pins, Low

The intended purpose of this BCLR instruction is to switch $P4^{0}$ from output to input.

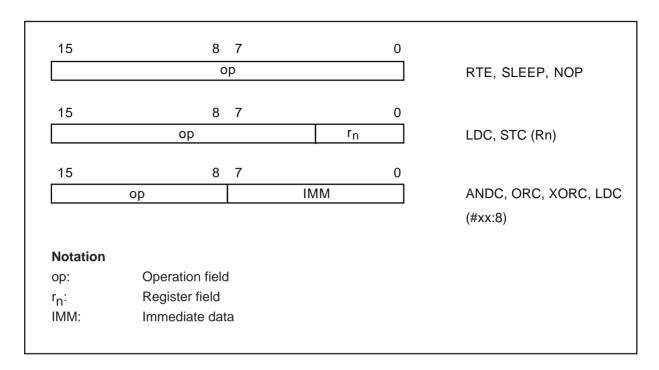


Figure 1-9. Machine-Language Coding of System Control Instructions

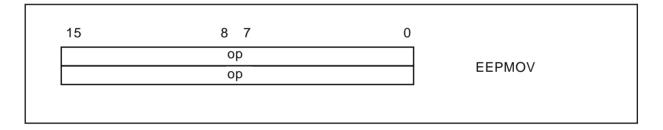


Figure 1-10. Machine-Language Coding of Block Data Transfer Instruction

1.3.4 Addressing Modes and Effective Address Calculation

Table 1-11 lists the eight addressing modes and their assembly-language notation. Each instruction can use a specific subset of these addressing modes.

Table 1-11. Addressing Modes

No.	Mode	Notation
(1)	Register direct	Rn
(2)	Register indirect	@Rn
(3)	Register indirect with 16-bit displacement	@(d:16, Rn)
(4)	Register indirect with post-increment	@Rn+
	Register indirect with pre-decrement	@-Rn
(5)	Absolute address (8 or 16 bits)	@aa:8, @aa:16
(6)	Immediate (3-, 8-, or 16-bit data)	#xx:3, #xx:8, #xx:16
(7)	PC-relative (8-bit displacement)	@(d:8, PC)
(8)	Memory indirect	@@aa:8

- (1) **Register Direct—Rn:** The register field of the instruction specifies an 8- or 16-bit general register containing the operand. In most cases the general register is accessed as an 8-bit register. Only the MOV.W, ADD.W, SUB.W, CMP.W, ADDS, SUBS, MULXU (8 bits \times 8 bits), and DIVXU (16 bits \div 8 bits) instructions have 16-bit operands.
- (2) **Register indirect**—@**Rn:** The register field of the instruction specifies a 16-bit general register containing the address of the operand.
- (3) Register Indirect with Displacement—@(d:16, Rn): This mode, which is used only in MOV instructions, is similar to register indirect but the instruction has a second word (bytes 3 and 4) which is added to the contents of the specified general register to obtain the operand address. For the MOV.W instruction, the resulting address must be even.

(4) Register Indirect with Post-Increment or Pre-Decrement—@Rn+ or @-Rn:

• Register indirect with post-increment—@Rn+
The @Rn+ mode is used with MOV instructions that load register from memory.
It is similar to the register indirect mode, but the 16-bit general register specified in the register field of the instruction is incremented after the operand is accessed. The size of the increment is 1 or 2 depending on the size of the operand: 1 for a byte operand; 2 for a

- word operand. For a word operand, the original contents of the 16-bit general register must be even.
- Register indirect with pre-decrement—@—Rn

 The @—Rn mode is used with MOV instructions that store registers contents to memory.

 It is similar to the register indirect mode, but the 16-bit general register specified in the register field of the instruction is decremented before the operand is accessed. The size of the decrement is 1 or 2 depending on the size of the operand: 1 for a byte operand; 2 for a word operand. For a word operand, the original contents of the 16-bit general register must be even.
- (5) **Absolute Address**—@aa:8 or @aa:16: The instruction specifies the absolute address of the operand in memory. The @aa:8 mode uses an 8-bit absolute address of the form H'FFxx. The upper 8 bits are assumed to be 1, so the possible address range is H'FF00 to H'FFFF (65280 to 65535). The MOV.B, MOV.W, JMP, and JSR instructions can use 16-bit absolute addresses.
- (6) Immediate—#xx:8 or #xx:16: The instruction contains an 8-bit operand in its second byte, or a 16-bit operand in its third and fourth bytes. Only MOV.W instructions can contain 16-bit immediate values.

The ADDS and SUBS instructions implicitly contain the value 1 or 2 as immediate data. Some bit manipulation instructions contain 3-bit immediate data (#xx:3) in the second or fourth byte of the instruction, specifying a bit number.

- (7) **PC-Relative**—@(**d:8, PC**): This mode is used to generate branch addresses in the Bcc and BSR instructions. An 8-bit value in byte 2 of the instruction code is added as a sign-extended value to the program counter contents. The result must be an even number. The possible branching range is -126 to +128 bytes (-63 to +64 words) from the current address.
- (8) Memory Indirect—@@aa:8: This mode can be used by the JMP and JSR instructions. The second byte of the instruction code specifies an 8-bit absolute address from H'0000 to H'00FF (0 to 255). Note that the initial part of the area from H'0000 to H'00FF contains the exception vector table. See the hardware manual of the specific chip for details. The word located at this address contains the branch address.

If an odd address is specified as a branch destination or as the operand address of a MOV.W instruction, the least significant bit is regarded as "0," causing word access to be performed at the address preceding the specified address. See the memory data structure description in section 1.1.2. Data Structure.

Calculation of Effective Address: Table 1-12 shows how the H8/300 calculates effective addresses in each addressing mode.

Arithmetic, logic, and shift instructions use register direct addressing (1). The ADD.B, ADDX, SUBX, CMP.B, AND, OR, and XOR instructions can also use immediate addressing (6).

The MOV instruction uses all the addressing modes except program-counter relative (7) and memory indirect (8).

Bit manipulation instructions use register direct (1), register indirect (2), or absolute (5) addressing to identify a byte operand and 3-bit immediate addressing to identify a bit within the byte. The BSET, BCLR, BNOT, and BTST instructions can also use register direct addressing (1) to identify the bit.

Effective Address Calculation

Table 1-12 explains how the effective address is calculated in each addressing mode.

Table 1-12, Effective Address Calculation (1)

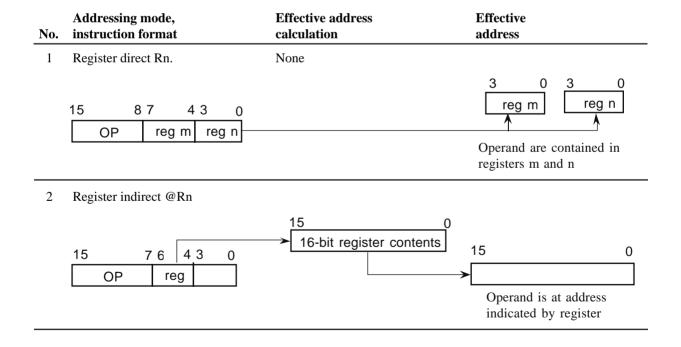
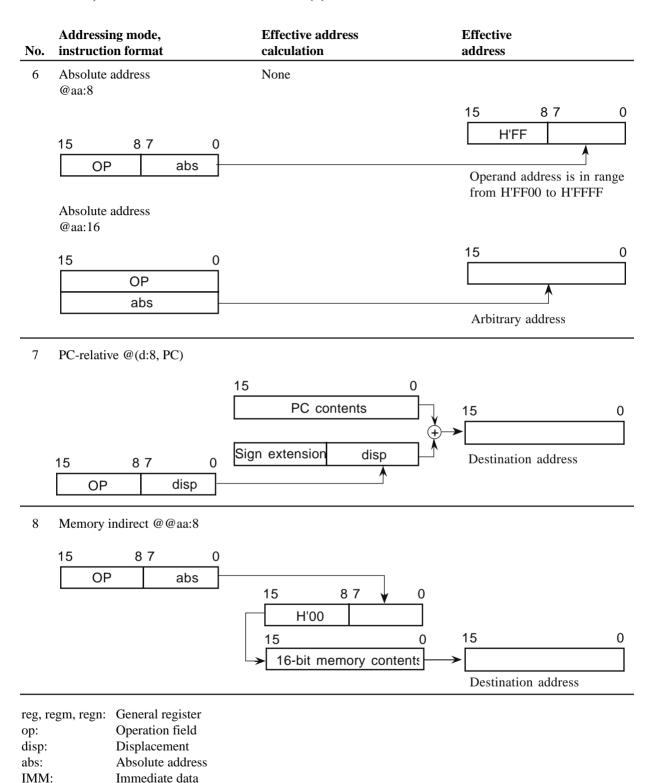


Table 1-12, Effective Address Calculation (2)

	•		
No.	Addressing mode, instruction format	Effective address calculation	Effective address
3	Register indirect with displacement $@(d:16,Rn)$		
	[15 0 16-bit register contents	15 0
	15 7 6 4 3 0 OP reg	16-bit displacement	Operand address is sum of register contents and displacement
4	Register indirect with pre-decremen @-Rn	t	
	15 7 6 4 3 0 OP reg	15 0 16-bit register contents 1 or 2*	Register is decremented before operand access
	Register indirect with post-increment @Rn+	nt	
		15 0 16-bit register contents for a byte operand, for a word operand	Register is incremented after operand access
5	Immediate #xx:8.	None	
	15 8 7 0 OP IMM		Operand is 1-byte immediate data
	Immediate #xx:16	None	
	15 0 OP IMM		Operand is 2-byte immediate data

Table 1-12, Effective Address Calculation (3)



Section 2. Instruction Set

Section 2 gives full descriptions of all the H8/300 instructions, presenting them in alphabetic order. Each instruction is explained in a table like the following:

ADD (ADD binary) (byte)

ADD

<Operation>

 $Rd + (EAs) \rightarrow Rd$

<Assembly-Language Format>

ADD.B <EAs>, Rd

<Examples>

ADD.B ROH, R1H

ADD.B #H'64, R2L

<Operand Size>

Byte

<Condition Code>

I	Η		N	Z	V	C
_	\$	_	\(\)	\$	‡	\$

- I: Previous value remains unchanged.
- H: Set to "1" when there is a carry from bit 3: otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Set to "1" if there is a carry from bit 7; otherwise cleared to "0."

<Description>

This instruction adds the source operand to the contents of an 8-bit general register and places the result in the 8-bit general register .

The source operand can be an 8-bit register value or immediate byte data.

Addressing	Mnem.	Operands		No. of					
mode	ivilletti.	Operands	1st byte		2nd byte		3rd byte	4th byte	states
Immediate	ADD.B	#xx:8, Rd	8	rd	IM	IM			2
Register direct	ADD.B	Rs, Rd	0	8	rs	rd			2

<Instruction Formats>

Name: The full and mnemonic names of the instruction are given at the top of the page.

Operation: Describes the instruction in symbolic notation. The following symbols are used.

Symbol	Meaning
(EAs)	Source operand
(EAd)	Destination operand
Rs, Rd, Rn	8-bit or 16-bit general register (s—source; d—destination)
#xx:3, #xx:8, #xx:16	3-bit, 8-bit, or 16-bit immediate data
d:8, d:16	8-bit or 16-bit displacement
PC	Program counter
SP	Stack pointer
CCR	Condition code register
Z	Zero flag in CCR
C	Carry flag in CCR
\rightarrow	The result of the operation on the left is assigned to the operand on the
	right (For compare instructions, the resulting condition code is
	assigned.)
+	Addition
_	Subtraction
×	Multiplication
÷	Division
^	AND logical
V	OR logical
\oplus	Exclusive OR logical
\leftrightarrow	Exchange
٦	Not

Assembly-Language

Format: The assemblylanguage coding of the $\frac{ADD.}{|}$ $\frac{B}{|}$ $\frac{\langle EAs \rangle}{|}$, $\frac{R}{|}$

instruction. An example is:

The operand size is indicated by the letter B (byte) or W (word). The size is indicated explicitly in this manual, but for instructions that permit only one size, the size designation can be omitted in source-program coding.

The abbreviation EAs or EAd (effective address of source or destination) is used for operands that permit more than one addressing mode.

Examples: Examples of the assembly-language coding of the instruction are given.

Operand size: Word or byte. Byte size is indicated for bit-manipulation instructions because these instructions access a full byte in order to read or write one bit.

Condition code: The effect of instruction execution on the flag bits in the CCR is indicated. The following notation is used:

Symbo	ol Meaning
‡	The flag is altered according to the result of the instruction.
0	The flag is cleared to "0."
	The flag is not changed.
*	Undetermined; the flag is left in an unpredictable state.

Description: A detailed explanation is given of the action of the instruction.

Instruction Formats: Each possible format of the instruction is shown explicitly, indicating the addressing mode, the object code, and the number of states required for execution when the instruction and its operands are located in on-chip memory. The following symbols are used:

Symbol	Meaning
Imm.	Immediate data (3, 8, or 16 bits)
abs.	An absolute address (8 bits or 16 bits)
disp.	Displacement (8 bits or 16 bits)
rs, rd, rn	General register number (3 bits or 4 bits) The s, d, and n correspond to the letters
	in the operand notation

16-bit general registers are indicated by a 3-bit rs, rd, or rn value. 8-bit registers are indicated by a 4-bit rs, rd, or rn value. Address registers used in the @Rn, @(disp:16, Rn), @Rn+, and @-Rn addressing modes are always 16-bit registers. Data registers are 8-bit or 16-bit registers depending on the size of the operand. For 8-bit registers, the lower three bits of rs, rd, or rn give the register number. The most significant bit is "1" if the lower byte of the register is used, or "0" if the upper byte is used. Registers are thus indicated as follows:

16-Bit register

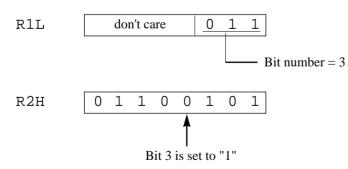
rs, rd, or rn		
Register		
0 0 0	R0	
0 0 1	R1	
:	:	
111	R7	

8-Bit registers

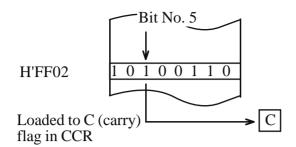
rs, rd, or rn	Register	
$0\ 0\ 0\ 0$	R0H	
0001	R1H	
:	:	
0 1 1 1	R7H	
1000	R0L	
1001	R1L	
:	:	
1111	R7L	

Bit Data Access: Bit data are accessed as the n-th bit of a byte operand in a general register or memory. The bit number is given by 3-bit immediate data, or by a value in a general register. When a bit number is specified in a general register, only the lower three bits of the register are significant. Two examples are shown below.





BLD #5, @H'FF02:8



The addressing mode and operand size apply to the register or memory byte containing the bit.

Number of States Required for Execution: The number of states indicated is the number required when the instruction and any memory operands are located in on-chip ROM or RAM. If the instruction or an operand is located in external memory or the on-chip register field, additional states are required for each access. See Appendix C.

 $Rd + (EAs) \rightarrow Rd$

<Assembly-Language Format>

ADD.B <EAs>, Rd

<Examples>

ADD.B ROH, R1H ADD.B #H'64, R2L

<Operand Size>

Byte

<Condition Code>

I	Η		N	Z	V	C
_	 \$	_		\$	\$	

- I: Previous value remains unchanged.
- H: Set to "1" when there is a carry from bit 3; otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Set to "1" if there is a carry from bit 7; otherwise cleared to "0."

<Description>

This instruction adds the source operand to the contents of an 8-bit general register and places the result in the 8-bit general register .

The source operand can be an 8-bit register value or immediate byte data.

Addressing	Mnem.	Operanda		No. of				
mode	iviriem.	Operands	1st byte		2nd byte	3rd byte	4th byte	states
Immediate	ADD.B	#xx:8, Rd	8	rd	IMM			2
Register direct	ADD.B	Rs, Rd	0	8	rs rd			2

 $Rd + Rs \rightarrow Rd$

<Assembly-Language Format>

ADD.W Rs, Rd

<Examples>

ADD.W R0, R1

<Operand Size>

Word

<Condition Code>

I	Н			N	Z	V	C
	_	‡			‡		\leftrightarrow

- I: Previous value remains unchanged.
- H: Set to "1" when there is a carry from bit 11; otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Set to "1" if there is a carry from bit 15; otherwise cleared to "0."

<Description>

This instruction adds word data in two general registers and places the result in the second general register.

Addressing	Mnem.	Operands		No. of					
mode	winem.	Operands	1st byte		2nd byte	3rd byte	4th byte	states	
Register direct	ADD.W	Rs, Rd	0	9	0 rs 0 rd			2	

 $Rd + 1 \longrightarrow Rd$

 $Rd + 2 \longrightarrow Rd$

<Assembly-Language Format>

ADDS #1, Rd

ADDS #2, Rd

<Examples>

ADDS #1, R4

ADDS #2, R5

<Operand Size>

Word

<Condition Code>

I	Η		N	Z	V	C
		_		_		_

I: Previous value remains unchanged.

H: Previous value remains unchanged.

N: Previous value remains unchanged.

Z: Previous value remains unchanged.

V: Previous value remains unchanged.

C: Previous value remains unchanged.

<Description>

This instruction adds the immediate value 1 or 2 to word data in a general register. Differing from the ADD instruction, it does not affect the condition code flags.

<Instruction Formats>

Addressing Mnem.		n. Operands		Instruction code							
mode	ivirietti.	Operands	1st b	1st byte 2nd byte		te	3rd byte	4th byte	states		
Register direct	ADDS	#1, Rd	0	В	0	0	rd			2	
Register direct	ADDS	#2, Rd	0	В	8	0	rd			2	

Note: This instruction cannot access byte size data.

 $Rd + (EAs) + C \longrightarrow Rd$

<Assembly-Language Format>

ADDX <EAs>, Rd

<Examples>

ADDX ROL, R1L ADDX #H'OA, R2H

<Operand Size>

Byte

<Condition Code>

I		Н	N	Z	V	C
	_		\Rightarrow	\leftrightarrow	\leftrightarrow	\Rightarrow

- I: Previous value remains unchanged.
- H: Set to "1" if there is a carry from bit 3; otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Set to "1" if there is a carry from bit 7; otherwise cleared to "0."

<Description>

This instruction adds the source operand and carry flag to the contents of an 8-bit general register and places the result in the 8-bit general register.

The source operand can be an 8-bit register value or immediate byte data.

Addressing Mnem.		Operands		Instruction code							
mode	iviriem.	Operands	1st b	yte	2nd	byte	3rd byte	4th byte	states		
Immediate	ADDX	#xx:8, Rd	9	rd	IIV	1M			2		
Register direct	ADDX	Rs, Rd	0	E	rs	rd			2		

 $Rd \wedge (EAs) \rightarrow Rd$

<Assembly-Language Format>

AND <EAs>, Rd

<Examples>

AND R6H, R6L AND #H'FD, R0H

<Operand Size>

Byte

<Condition Code>

I		Н		N	Z	V	C
_	_	_	_		‡	0	

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction ANDs the source operand with the contents of an 8-bit general register and places the result in the 8-bit general register.

The source operand can be an 8-bit register value or immediate byte data.

Addressing	Mnem.	Operanda		No. of						
mode	willem.	Operands	1st b	1st byte 2nd		byte	3rd byte	4th byte	states	
Immediate	AND	#xx:8, Rd	Е	rd	IIV	IM			2	
Register direct	AND	Rs, Rd	1	6	rs	rd			2	

ANDC (AND Control register)

ANDC

<opera< th=""><th>ation></th></opera<>	ation>
---	--------

 $CCR \wedge \#IMM \rightarrow CCR$

<Assembly-Language Format>

ANDC #xx:8, CCR

<Examples>

ANDC #H'7F, CCR

<Operand Size>

Byte

<Condition Code>

I		Η		N	Z	V	C
‡		\$			\$	\$	

- I: ANDed with bit 7 of the immediate data.
- H: ANDed with bit 5 of the immediate data.
- N: ANDed with bit 3 of the immediate data.
- Z: ANDed with bit 2 of the immediate data.
- V: ANDed with bit 1 of the immediate data.
- C: ANDed with bit 0 of the immediate data.

<Description>

This instruction ANDs the condition code register (CCR) with immediate data and places the result in the condition code register. Bits 6 and 4 are ANDed as well as the flag bits.

No interrupt requests are accepted immediately after this instruction. All interrupts, including the nonmaskable interrupt (NMI), are deferred until after the next instruction.

Addressing	Mnem. Operands -			No. of				
mode	Willetti.	Operanus	1st byte		2nd byte	3rd byte	4th byte	states
Immediate	ANDC	#xx:8, CCR	0	6	IMM			2

 $C \wedge (<Bit No.> of <EAd>) \rightarrow C$

<Assembly-Language Format>

BAND #xx:3, <EAd>

<Examples>

BAND #0, R1L

BAND #4, @R3

BAND #7, @H'FFE0:8

<Condition Code>

I		Η		N	Z	V	C
	_		_			_	

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: ANDed with the specified bit.

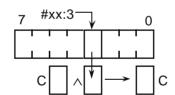
<Operand Size>

Byte

<Description>

This instruction ANDs a specified bit with the carry flag and places the result in the carry flag. The specified bit can be located in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

 $\label{eq:energy} \hbox{Bit No.}$ $<\!\!EAd\!\!>^*\!\!\to\!Byt\!e$ data in register or memory



The value of the specified bit is not changed.

Addressing	Mnem. Operands			Instruction code							
mode	wiriem.	Operands	1st b	oyte	2nd b	yte	3rd b	yte	4th l	oyte	states
Register direct	BAND	#xx:3, Rd	7	6	0 IMM	rd					2
Register indirect	BAND	#xx:3,@Rd	7	С	0 rd	0	7	6	0 IMM	0	6
Absolute address	BAND	#xx:3,@aa:8	7	E	al	os	7	6	0 IMM	0	6

^{*} Register direct, register indirect, or absolute addressing.

If cc then

$$PC + d:8 \rightarrow PC$$

else next;

<Assembly-Language Format>

B<u>c</u>c d:8

Condition code field

(For mnemonics, see the table on the next page.)

<Examples>

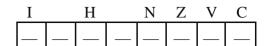
BHI H'42

BEQ H'-7E

<Operand Size>

—

<Condition Code>



- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

If the specified condition is false, this instruction does nothing; the next instruction is executed. If the specified condition is true, a signed displacement is added to the address of the next instruction and execution branches to the resulting address.

The displacement is a signed 8-bit value which must be even. The branch destination address can be located in the range -126 to +128 bytes from the address of the Bcc instruction.

The available conditions and their mnemonics are given below.

Mnemonic	cc Field	Description	Condition	Meaning
BRA (BT)	0000	Always (True)	Always true	
BRN (BF)	0001	Never (False)	Never (False) Never	
BHI	0010	HIgh	$C \bigvee Z = 0$	X > Y (Unsigned)
BLS	0011	Low or Same	$C_{\vee}Z=1$	$X \le Y$ (Unsigned)
BCC (BHS)	0100	Carry Clear (High or Same)	C = 0	$X \ge Y$ (Unsigned)
BCS (BLO)	0101	Carry Set (LOw)	C = 1	X < Y (Unsigned)
BNE	0110	Not Equal	Z = 0	$X \neq Y$ (Signed or unsigned)
BEQ	0111	EQual	Z = 1	X = Y (Signed or unsigned)
BVC	1000	oVerflow Clear	V = 0	
BVS	1001	oVerflow Set	V = 1	
BPL	1010	PLus	N = 0	
BMI	1011	MInus	N = 1	
BGE	1100	Greater or Equal	$N \oplus V = 0$	$X \ge Y$ (Signed)
BLT	1101	Less Than	$N \oplus V = 1$	X < Y (Signed)
BGT	1110	Greater Than	$Z_{\vee}(N \oplus V) = 0$	X > Y (Signed)
BLE	1111	Less or Equal	$Z_{\vee}(N \oplus V) = 1$	$X \le Y$ (Signed)

BT, BF, BHS, and BLO are synonyms for BRA, BRN, BCC, and BCS, respectively.

Adressing					Instruction	on code		No . of
mode	Mnem.	Operands	1st byte		2nd byte	3rd byte	4th byte	states
PC relative	BRA (BT)	d:8	4	0	disp.			4
PC relative	BRN (BF)	d:8	4	1	disp.			4
PC relative	ВНІ	d:8	4	2	disp.			4
PC relative	BLS	d:8	4	3	disp.			4
PC relative	BCC (BHS)	d:8	4	4	disp.			4
PC relative	BCS (BLO)	d:8	4	5	disp.			4
PC relative	BNE	d:8	4	6	disp.			4
PC relative	BEQ	d:8	4	7	disp.			4
PC relative	BVC	d:8	4	8	disp.			4
PC relative	BVS	d:8	4	9	disp.			4
PC relative	BPL	d:8	4	А	disp.			4
PC relative	BMI	d:8	4	В	disp.			4
PC relative	BGE	d:8	4	С	disp.			4
PC relative	BLT	d:8	4	D	disp.			4
PC relative	BGT	d:8	4	Е	disp.			4
PC relative	BLE	d:8	4	F	disp.			4

^{*} The branch address must be even.

 $0 \longrightarrow (<Bit No.> of <EAd>)$

<Assembly-Language Format>

BCLR #xx:3, <EAd>

<Examples>

BCLR #0, ROL

BCLR #1, @R5

BCLR R6L, @H'FFCO:8

<Operand Size>

Byte

<Condition Code>

I		Н		N	Z	V	C
_	_		_	_	_		_

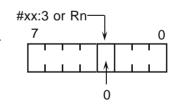
- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction clears a specified bit in the destination operand to "0." The bit number can be specified by 3-bit immediate data, or by the lower three bits of an 8-bit general register. The destination operand can be located in a general register or memory.

The specified bit is not tested before being cleared. The condition code flags are not altered.

 $\label{eq:energy} \hbox{Bit No.} $$<EAd>* \to Byte data in register or memory$



^{*}Register direct, register indirect, or absolute addressing.

BCLR (Bit CLeaR) BCLR

Addressing	Mnom	Operanda		Ins	truction	n code						No. of
mode	Mnem.	Operands	1st b	yte	2nd	byte	3rd b	yte	41	h k	oyte	states
Register direct	BCLR	#xx:3, Rd	7	2	0 IMM	rd						2
Register indirect	BCLR	#xx:3,@Rd	7	D	0 rd	0	7	2	o¦ IN	ΙM	0	8
Absolute address	BCLR	#xx:3,@aa:8	7	F	al	bs	7	2	0 IN	ΙM	0	8
Register direct	BCLR	Rn, Rd	6	2	rn	rd						2
Register indirect	BCLR	Rn, @Rd	7	D	0 rd	0	6	2	rn	1	0	8
Absolute address	BCLR	Rn, @aa:8	7	F	ab	s	6	2	rn	-	0	8

 $C \cap (\langle Bit \text{ No.} \rangle \text{ of } \langle EAd \rangle)] \rightarrow C$

<Assembly-Language Format>

BIAND #xx:3, <EAd>

<Examples>

BIAND #0, R1H

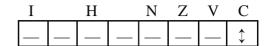
BIAND #2, @R5

BIAND #4, @H'FFDE:8

<Operand Size>

Byte

<Condition Code>

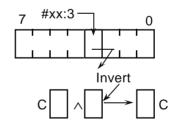


- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: ANDed with the inverse of the specified bit.

<Description>

This instruction ANDs the inverse of a specified bit with the carry flag and places the result in the carry flag. The specified bit can be located in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

 $\label{eq:bit No.} \mbox{\ensuremath{\mathsf{EAd}}}\mbox{\ensuremath{\mathsf{*}}} \mbox{\ensuremath{\mathsf{Byte}}} \mbox{\ensuremath{\mathsf{data}}} \mbox{\ensuremath{\mathsf{in}}} \mbox{\ensuremath{\mathsf{register}}} \mbox{\ensuremath{\mathsf{e}}} \mbox{\ens$



The value of the specified bit is not changed.

Addressing	Mnom	Inem. Operands -		Ins	tructio	n code					No. of
mode	wiriem.	Operanus	1st b	yte	2nd l	byte	3rd b	oyte	4th b	oyte	states
Register direct	BIAND	#xx:3, Rd	7	6	1 IMM	rd					2
Register indirect	BIAND	#xx:3,@Rd	7	С	0 rd	0	7	6	1 IMM	0	6
Absolute address	BIAND	#xx:3,@aa:8	7	E	a	bs	7	6	1 IMM	0	6

^{*}Register direct, register indirect, or absolute addressing.

 $\neg (\langle Bit \ No. \rangle \ of \langle EAd \rangle) \longrightarrow C$

<Assembly-Language Format>

BILD #xx:3, <EAd>

<Examples>

BILD #3, R4L

BILD #5, @R5

BILD #7, @H'FFA2:8

<Operand Size>

Byte

<Condition Code>

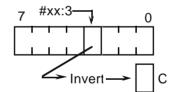
I	Η		N	Z	V	C
	 _	_	_	_		

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Loaded with the inverse of the specified bit.

<Description>

This instruction loads the inverse of a specified bit into the carry flag. The specified bit can be located in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

Bit No. $\langle EAd \rangle^* \rightarrow Byte data in register or memory$



The value of the specified bit is not changed.

Addressing	Mnem.	Operands	Ins	struction code			No. of
mode	wiriem.	Operanus	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	BILD	#xx:3, Rd	7 7	1 IMM rd			2
Register indirect	BILD	#xx:3,@Rd	7 C	0 rd 0	7 7	1 IMM 0	6
Absolute address	BILD	#xx:3,@aa:8	7 E	abs	7 7	1 IMM 0	6

^{*}Register direct, register indirect, or absolute addressing.

 $C \setminus [\neg (\langle Bit \ No. \rangle \ of \langle EAd \rangle)] \rightarrow C$

<Assembly-Language Format>

BIOR #xx:3, <EAd>

<Examples>

BIOR #6, R1H

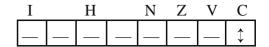
BIOR #3, @R2

BIOR #0, @H'FFF0:8

<Operand Size>

Byte

<Condition Code>

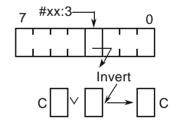


- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: ORed with the inverse of the specified bit.

<Description>

This instruction ORs the inverse of a specified bit with the carry flag and places the result in the carry flag. The specified bit can be located in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

 $\label{eq:energy} \hbox{Bit No.} $$<EAd>* \to Byte data in register or memory$



The value of the specified bit is not changed.

Addressing	Mnem.	Operands			No. of							
mode	wiriem.	Operands	1st byte		2nd byte		3rd l	oyte	4	th by	te	states
Register direct	BIOR	#xx:3, Rd	7	4	1 IMM	rd						2
Register indirect	BIOR	#xx:3,@Rd	7	С	0 rd	0	7	4	1 1 1	ИМ	0	6
Absolute address	BIOR	#xx:3,@aa:8	7	E	al	os	7	4	1 1 1	лм	0	6

^{*}Register direct, register indirect, or absolute addressing.

 $\neg C \rightarrow (\langle Bit \ No. \rangle \ of \langle EAd \rangle)$

<Assembly-Language Format>

BIST #xx:3, <EAd>

<Examples>

BIST #0, ROL

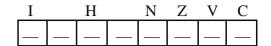
BIST #6, @R3

BIST #7, @H'FFBB:8

<Operand Size>

Byte

<Condition Code>

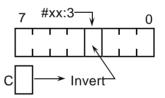


- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction stores the inverse of the carry flag to a specified bit location in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

 $\label{eq:energy} \hbox{Bit No.} $$<EAd>* \to Byte data in register or memory$



The values of the unspecified bits are not changed.

Addressing	Mnem. Operands			Instruction code								
mode	winem.	Operands	1st b	oyte	2nd b	oyte	3rd b	yte	4th b	yte	states	
Register direct	BIST	#xx:3, Rd	6	7	1 IMM	rd					2	
Register indirect	BIST	#xx:3,@Rd	7	D	0 rd	0	6	7	1 IMM	0	8	
Absolute address	BIST	#xx:3,@aa:8	7	F	al	os	6	7	1 IMM	0	8	

^{*} Register direct, register indirect, or absolute addressing.

 $C \oplus [\neg (\langle Bit \ No. \rangle \ of \langle EAd \rangle)] \xrightarrow{} C$

<Assembly-Language Format>

BIXOR #xx:3, <EAd>

<Examples>

BIXOR #1, R4L

BIXOR #2, @R5

BIXOR #3, @H'FF60:8

SIAOR #I, RAL

Z: Previous value remains unchanged.

<Condition Code>

I:

N:

V: Previous value remains unchanged.

C: Exclusive-ORed with the inverse of the specified bit.

Previous value remains unchanged.

Previous value remains unchanged.

Previous value remains unchanged.

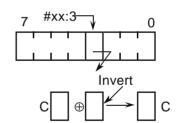
<Operand Size>

Byte

<Description>

This instruction exclusive-ORs the inverse of a specified bitwith the carry flag and places the result in the carry flag. The specified bit can be located in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

 $\label{eq:energy} \hbox{Bit No.} $$<EAd>*\rightarrow Byte data in register or memory$



The value of the specified bit is not changed.

Addressing	Mnem.	Operanda		Ins	struction	code					No. of
mode	winem.	Operands	1st l	oyte	2nd b	yte	3rd l	oyte	4th	byte	states
Register direct	BIXOR	#xx:3, Rd	7	5	1 IMM	rd					2
Register indirect	BIXOR	#xx:3,@Rd	7	С	0 rd	0	7	5	1 IMM	0	6
Absolute address	BIXOR	#xx:3,@aa:8	7	E	al)S	7	5	1¦IMM	0	6

^{*} Register direct, register indirect, or absolute addressing.

BLD (Bit LoaD)

BLD

<Operation>

(<Bit No.> of <EAd $>) _ C$

<Assembly-Language Format>

BLD #xx:3, <EAd>

<Examples>

BLD #1, R3H

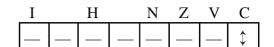
BLD #2, @R2

BLD #4, @H'FF90:8

<Operand Size>

Byte

<Condition Code>

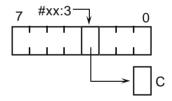


- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Loaded with the specified bit.

<Description>

This instruction loads a specified bit into the carry flag. The specified bit can be located in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

 $\label{eq:bit No.} $$ \ensuremath{^{<\!}EAd>^* \to Byte data in register or memory} $$$



The value of the specified bit is not changed.

Addressing	Mnem.	Operands		Ins	struction	ode					No. of
mode	wiriem.	Operands	1st l	byte	2nd b	yte	3rd b	oyte	4th b	oyte	states
Register direct	BLD	#xx:3, Rd	7	7	0 IMM	rd					2
Register indirect	BLD	#xx:3,@Rd	7	С	0 rd	0	7	7	0 IMM	0	6
Absolute address	BLD	#xx:3,@aa:8	7	ļΕ	al	os	7	7	0 IMM	0	6

^{*} Register direct, register indirect, or absolute addressing.

¬ (<Bit No.> of <EAd>)
__ (<Bit No.> of <EAd>)

<Assembly-Language Format>

BNOT #xx:3, <EAd>

BNOT Rn, <EAd>

<Examples>

BNOT #7, R1H

BNOT R1L, @R6

BNOT #3, @H'FFB4:8

<Operand Size>

Byte

<Condition Code>

I	Н	N	Z	V	C
	 			_	

I: Previous value remains unchanged.

H: Previous value remains unchanged.

N: Previous value remains unchanged.

Z: Previous value remains unchanged.

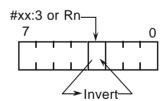
V: Previous value remains unchanged.

C: Previous value remains unchanged.

<Description>

This instruction inverts a specified bit in a general register or memory location. The bit number is specified by 3-bit immediate data, or by the lower three-bits of a general register. The operation is shown schematically below.

 $\hbox{ Bit No. } <\!\!EAd>^*\!\!\rightarrow\! Byte data in register or memory$



The bit is not tested before being inverted. The condition code flags are not altered.

^{*}Register direct, register indirect, or absolute addressing.

BNOT (Bit NOT)

BNOT

Addressing	Managa	On a ran da		Instruction code									
mode	Mnem.	Operands	1st b	yte	2nd	byte	3rd b	yte	4th	byte	states		
Register direct	BNOT	#xx:3, Rd	7	1	0 IMM	rd					2		
Register indirect	BNOT	#xx:3,@Rd	7	D	0 rd	0	7	1	0 IMM	0	8		
Absolute address	BNOT	#xx:3,@aa:8	7	F	а	bs	7	1	0 IMM	0	8		
Register direct	BNOT	Rn, Rd	6	1	rn	rd					2		
Register indirect	BNOT	Rn, @Rd	7	D	0 rd	0	6	1	rn	0	8		
Absolute address	BNOT	Rn, @aa:8	7	F	at	os	6	1	rn	0	8		

 $C \lor (<Bit No.> of <EAd>) \rightarrow C$

<Assembly-Language Format>

BOR #xx:3, <EAd>

<Examples>

BOR #5, R2H

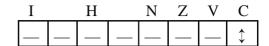
BOR #4, @R1

BOR #5, @H'FFB6:8

<Operand Size>

Byte

<Condition Code>

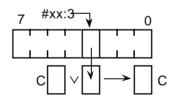


- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: ORed with the specified bit.

<Description>

This instruction ORs a specified bit with the carry flag and places the result in the carry flag. The specified bit can be located in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

 $\label{eq:energy} \hbox{Bit No.} $$<EAd>* \to Byte data in register or memory$



The value of the specified bit is not changed.

^{*}Register direct, register indirect, or absolute addressing.

Addressing	Mnem.	Operanda		No. of								
mode	wiriem.	Operands	1st byte		2nd byte		3rd byte		4th byte		states	
Register direct	BOR	#xx:3, Rd	7	4	0 IMM	rd					2	
Register indirect	BOR	#xx:3,@Rd	7	С	0 rd	0	7	4	0 IMM	0	6	
Absolute address	BOR	#xx:3,@aa:8	7	E	al	os	7	4	0 IMM	0	6	

 $1 \longrightarrow (<Bit No.> of <EAd>)$

<Assembly-Language Format>

BSET #xx:3,<EAd>

BSET Rn,<EAd>

<Examples>

BSET #3, R2L

BSET R2H, @R6

BSET #7, @H'FFE4:8

<Operand Size>

Byte

<Condition Code>

I	Η		N	Z	V	C
		_	_	_		_

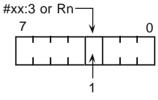
- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction sets a specified bit in the destination operand to "1." The bit number can be specified by 3-bit immediate data, or by the lower three-bits of an 8-bit general register. The destination operand can be located in a general register or memory.

The specified bit is not tested before being cleared. The condition code flags are not altered.

Bit No. <EAd>*→Byte data in register or memory



*Register direct, register indirect, or absolute addressing.

BSET (Bit SET)
BSET

Addressing mode	Maam	Operands			No. of						
	Mnem.		1st byte		2nd byte		3rd byte		4th byte		states
Register direct	BSET	#xx:3, Rd	7	0	0 IMM	rd					2
Register indirect	BSET	#xx:3,@Rd	7	D	0 rd	0	7	0	0 IMM	0	8
Absolute address	BSET	#xx:3,@aa:8	7	F	al	os	7	0	0 IMM	0	8
Register direct	BSET	Rn, Rd	6	0	rn	rd					2
Register indirect	BSET	Rn, @Rd	7	D	0 rd	0	6	0	rn	0	8
Absolute address	BSET	Rn, @aa:8	7	F	ab	S	6	0	rn	0	8

 $PC \rightarrow @-SP$

 $PC + d:8 \rightarrow PC$

<Condition Code>

I H N Z V C

<Assembly-Language Format>

BSR d:8

<Examples>

BSR H'76

<Operand Size>

I: Previous value remains unchanged.

H: Previous value remains unchanged.

N: Previous value remains unchanged.

Z: Previous value remains unchanged.

V: Previous value remains unchanged.

C: Previous value remains unchanged.

<Description>

This instruction pushes the program counter (PC) value onto the stack, then adds a specified displacement to the program counter value and branches to the resulting address. The program counter value used is the address of the instruction following the BSR instruction.

The displacement is a signed 8-bit value which must be even. The possible branching range is -126 to +128 bytes from the address of the BSR instruction.

Addressing mode	Mnem.	Operands		No. of				
	wiffern.	Operanus	1st b	oyte 2nd byte		nd byte 3rd byte 4th byte		states
PC-relative	BSR	d:8	5	5	disp			6

 $C \rightarrow (<Bit No.> of <EAd>)$

<Assembly-Language Format>

BST #xx:3, <EAd>

<Examples>

BST #7, R4L

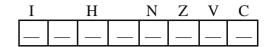
BST #2, @R3

BST #6, @H'FFD1:8

<Operand Size>

Byte

<Condition Code>



I: Previous value remains unchanged.

H: Previous value remains unchanged.

N: Previous value remains unchanged.

Z: Previous value remains unchanged.

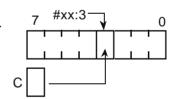
V: Previous value remains unchanged.

C: Previous value remains unchanged.

<Description>

This instruction stores the carry flag to a specified flag location in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

 $\hbox{$$ \hbox{Bit No.}$} < \hbox{$EAd$} > * \rightarrow \hbox{$Byte data in register or memory}$



Addressing mode	Mnem.	Operando		No. of						
	ivilletti.	Operands	1st l	byte	2nd b	oyte	3rd b	oyte	4th byte	states
Register direct	BST	#xx:3, Rd	6	7	0 IMM	rd				2
Register indirect	BST	#xx:3,@Rd	7	D	0 rd	0	6	7	0 IMM 0	8
Absolute address	BST	#xx:3,@aa:8	7	F	al	os	6	7	0 IMM 0	8

^{*} Register direct, register indirect, or absolute addressing.

 $\neg (\langle Bit \ No. \rangle \ of \langle EAd \rangle) \xrightarrow{} Z$

<Assembly-Language Format>

BTST #xx:3, <EAd>

BTST Rn, <EAd>

<Examples>

BTST #4, R6L

BTST R1H, @R5

BTST #7, @H'FF6C:8

<Operand Size>

Byte

<Condition Code>

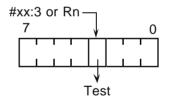
I		Н		N	Z	V	C
	_		_		\$	_	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: See to "1" if the specified bit is zero; otherwise cleared to "0".
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction tests a specified bit in a general register or memory location and sets or clears the Zero flag accordingly. The bit number can be specified by 3-bit immediate data, or by the lower three bits of an 8-bit general register. The operation is shown schematically below.

 $\label{eq:bit No.} \mathsf{EAd} \mathsf{>} \mathsf{*} \!\!\to\! \mathsf{Byte} \ \mathsf{data} \ \mathsf{in} \ \mathsf{register} \ \mathsf{or} \ \mathsf{memory}$



The value of the specified bit is not altered.

^{*}Register direct, register indirect, or absolute addressing.

BTST (Bit TeST)
BTST

Addressing	Mnem.	Operanda			No. of						
mode	winem.	Operands	1st byte		2nd byte	3rd byte		4th byte		states	
Register direct	BTST	#xx:3, Rd	7	3	0 IMM rd					2	
Register indirect	BTST	#xx:3,@Rd	7	С	0 rd 0	7	3	0 IMM	0	6	
Absolute address	BTST	#xx:3,@aa:8	7	Е	abs	7	3	0 IMM	0	6	
Register direct	BTST	Rn, Rd	6	3	rn rd					2	
Register indirect	BTST	Rn, @Rd	7	С	0 rd 0	6	3	rn	0	6	
Absolute address	BTST	Rn, @aa:8	7	E	abs	6	3	rn	0	6	

 $C \oplus (\langle Bit \text{ No.} \rangle \text{ of } \langle EAd \rangle) \rightarrow C$

<Assembly-Language Format>

BXOR #xx:3, <EAd>

<Examples>

BXOR #4, R6H

BXOR #2, @R0

BXOR #1, @H'FFA0:8

<Operand Size>

Byte

<Condition Code>

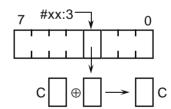
Ι	Н		N	Z	V	C
_	 _	_	_	_		

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Exclusive-ORed with the specified bit.

<Description>

This instruction exclusive-ORs a specified bit with the carry flag and places the result in the carry flag. The specified bit can be located in a general register or memory. The bit number is specified by 3-bit immediate data. The operation is shown schematically below.

Bit No. $\langle EAd \rangle^* \rightarrow Byte data in register or memory$



The value of the specified bit is not changed.

^{*}Register direct, register indirect, or absolute addressing.

Addressing Mnem.		Operanda		Instruction code								No. of
mode	winem.	Operands	1st b	oyte	2nd b	yte	3rd b	yte	41	h b	yte	states
Register direct	BXOR	#xx:3, Rd	7	5	0 IMM	rd						2
Register indirect	BXOR	#xx:3,@Rd	7	С	0 rd	0	7	5	0¦IN	/M	0	6
Absolute address	BXOR	#xx:3,@aa:8	7	¦Ε	al	os	7	5	0¦IN	/M	0	6

Rd – (EAs); set condition code

<Assembly-Language Format>

CMP.B <EAs>, Rd

<Examples>

CMP.B #H'E5, R1H CMP.B R3L, R4L

<Operand Size>

Byte

<Condition Code>

I		Н	N	Z	V	С
_	_	\$	\(\)	\$	\leftrightarrow	\$

- I: Previous value remains unchanged.
- H: Set to "1" when there is a borrow from bit 3; otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Set to "1" if there is a borrow from bit 7; otherwise cleared to "0."

<Description>

This instruction subtracts an 8-bit source register or immediate data from an 8-bit destination register and sets the condition code flags according to the result. The destination register is not altered.

Addressing	Mnem.	Operanda			No. of				
mode	winem.	Operands	1st byte		2nd byte		3rd byte	4th byte	states
Immediate	CMP.B	#xx:8,Rd	Α	rd	IN	1M			2
Register direct	CMP.B	Rs, Rd	1	С	rs	rd			2

Rd - Rs; set condition code

<Assembly-Language Format>

CMP.W Rs, Rd

<Examples>

CMP.W R5, R6

<Operand Size>

Word

<Condition Code>

I		Н		N	Z	V	C
	_	\$	_				1

- I: Previous value remains unchanged.
- H: Set to "1" when there is a borrow from bit 11; otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Set to "1" if there is a borrow from bit 15; otherwise cleared to "0."

<Description>

This instruction subtracts a source register from a destination register and sets the condition code flags according to the result. The destination register is not altered.

Addressing	Mnem. Operands			No. of			
mode	iviriem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	CMP.W	Rs, Rd	1 D	0 rs 0 rd			2

Rd (decimal adjust) → Rd

<Assembly-Language Format>

DAA Rd

<Examples>

DAA R5L

<Operand Size>

Byte

<Condition Code>

I		Н	N	Z	V	C
	_	*	\Rightarrow	‡	*	\leftrightarrow

- I: Previous value remains unchanged.
- H: Unpredictable.
- N: Set to "1" if the adjusted result is negative; otherwise cleared to "0."
- Z: Set to "1" if the adjusted result is zero; otherwise cleared to "0."
- V: Unpredictable.
- C: Set to "1" if there is a carry from bit 7; otherwise left unchanged.

<Description>

Given that the result of an addition operation performed by the ADD.B or ADDX instruction on 4-bit BCD data is contained in an 8-bit general register and the carry and half-carry flags, the DAA instruction adjusts the result by adding H'00, H'06, H'60, or H'66 to the general register according to the table below.

Valid results are not assured if this instruction is executed under conditions other than those stated above.

	Status before	e adjustme	nt	Value	Resulting
C flag	Upper nibble	H flag	Lower nibble	added	C flag
0	0 – 9	0	0 – 9	H'00	0
0	0 - 8	0	A - F	H'06	0
0	0 – 9	1	0 - 3	H'06	0
0	A - F	0	0 – 9	H'60	1
0	9 – F	0	A - F	H'66	1
0	A - F	1	0 - 3	H'66	1
1	0 - 2	0	0 – 9	H'60	1
1	0 - 2	0	A - F	H'66	1
1	0 - 3	1	0 - 3	H'66	1

Addressing	Mnom	Operands			In	structi	on code		No. of
mode	Mnem.	Operands	1st byte		2nd byte		3rd byte	4th byte	states
Register direct	DAA	Rd	0	F	0	rd			2

Rd (decimal adjust) → Rd

<Assembly-Language Format>

DAS Rd

<Examples>

DAS ROH

<Operand Size>

Byte

<Condition Code>

I	Η	N	Z	V	С
	 *	‡	\$	*	

- I: Previous value remains unchanged.
- H: Unpredictable.
- N: Set to "1" if the adjusted result is negative; otherwise cleared to "0."
- Z: Set to "1" if the adjusted result is zero; otherwise cleared to "0."
- V: Unpredictable.
- C: Previous value remains unchanged.

<Description>

Given that the result of a subtraction operation performed by the SUB.B, SUBX, or NEG instruction on 4-bit BCD data is contained in an 8-bit general register and the carry and half-carry flags, the DAA instruction adjusts the result by adding H'00, H'FA, H'A0, or H'9A to the general register according to the table below.

Valid results are not assured if this instruction is executed under conditions other than those stated above.

	Status before	e adjustme	nt	Value	Resulting
C flag	Upper nibble	H flag	Lower nibble	added	C flag
0	0 – 9	0	0 – 9	H'00	0
0	0 - 8	1	6 – F	H'FA	0
1	7 – F	0	0 – 9	H'A0	1
1	6 – F	1	6 – F	H'9A	1

Addressing	Mnem. Operands				No. of		
mode	iviriem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	DAS	Rd	1 F	0 rd			2

 $Rd - 1 \longrightarrow Rd$

<Assembly-Language Format>

DEC Rd

<Examples>

DEC R2L

<Operand Size>

Byte

<Condition Code>

I	Н	N	Z	V	C
			\$		_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs (the previous value in Rd was H'80); otherwise cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction decrements an 8-bit general register and places the result in the 8-bit general register.

Addressing	Mnem.	Operands		No. of			
mode	winem.		1st byte	2nd byte	3rd byte	4th byte	states
Register direct	DEC	Rd	1 A	0 rd			2

 $Rd \div Rs \longrightarrow Rd$

<Assembly-Language Format>

DIVXU Rs, Rd

<Examples>

DIVXU ROL, R1

<Operand Size>

Byte

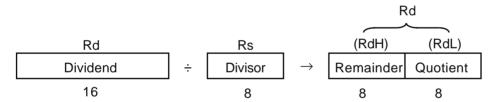
<Condition Code>

I		Η	N	Z	V	C
_	_		\(\)	\(\)		_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the divisor is negative; otherwise cleared to "0."
- Z: Set to "1" if the divisor is zero; otherwise cleared to "0."
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction divides a 16-bit general register by an 8-bit general register and places the result in the 16-bit general register. The quotient is placed in the lower byte. The remainder is placed in the upper byte. The operation is shown schematically below.



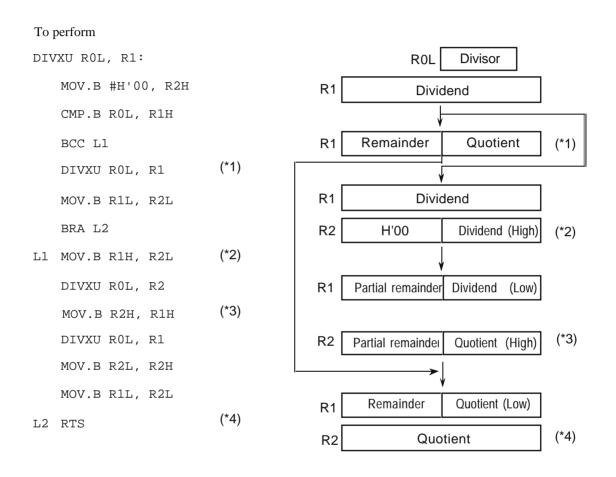
Valid results are not assured if division by zero is attempted or an overflow occurs. Division by zero is indicated in the Zero flag. Overflow can be avoided by the coding shown on the next page.

Addressing	Mnom	Mnem. Operands		Instruction code						
mode	iviriem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states			
Register direct	DIVXU	Rs, Rd	5 1	rs 0 rd			14			

<Note: DIVXU Overflow>

Since the DIVXU instruction performs 16-bit \div 8-bit \longrightarrow 8-bit division, an overflow will occur if the divisor byte is equal to or less than the upper byte of the dividend. For example, H'FFFF \div H'01 \longrightarrow H'FFFF causes an overflow. (The quotient has more than 8 bits.)

Overflows can be avoided by using a subprogram like the following. A work register is required.



if R4L \neq 0 then repeat @R5+ \rightarrow @R6+ R4L - 1 \rightarrow R4L until R4L = 0

else next;

<Assembly-Language Format>

EEPMOV

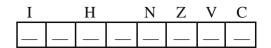
<Examples>

MOV.B #H'20, R4L MOV.W #H'FEC0, R5 MOV.W #H'6000, R6 EEPMOV

<Operand Size>

_

<Condition Code>



- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction moves a block of data from the memory location specified in general register R5 to the memory location specified in general register R6. General register R4L gives the byte length of the block.

Data are transferred a byte at a time. After each byte transfer, R5 and R6 are incremented and R4L is decremented. When R4L reaches 0, the transfer ends and the next instruction is executed. No interrupt requests are accepted during the data transfer.

At the end of this instruction, R4L contains H'00. R5 and R6 contain the last transfer address +1.

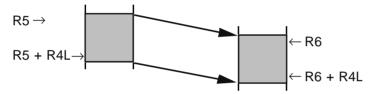
Chips in the H8/300 Series having large on-chip EEPROM memories use this instruction to write data in the EEPROM. For details, see the hardware manual for the particular chip. The memory locations specified by general registers R5 and R6 are read before the block transfer is performed.

<Instruction Formats>

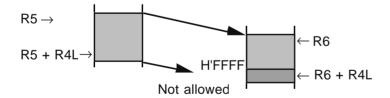
Addressing	Mnem. Operands		Inst	No. of			
mode	winem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
_	EEPMOV		7 B	5 C	5 9	8 F	8+4n*

Notes on EEPMOV Instruction

1. The EEPMOV instruction is a block data transfer instruction. It moves the number of bytes specified by R4L from the address specified by R5 to the address specified by R6.



2. When setting R4L and R6, make sure that the final destination address (R6 + R4L) does not exceed H'FFFF. The value in R6 must not change from H'FFFF to H'0000 during execution of the instruction.



^{*} n is the initial value in R4L ($0 \le n \le 255$). Although n bytes of data are transferred, memory is accessed 2(n+1) times, requiring 4(n+1) states.

 $Rd + 1 \longrightarrow Rd$

<Assembly-Language Format>

INC Rd

<Examples>

INC R3L

<Operand Size>

Byte

<Condition Code>

I		Н		N	Z	V	C
	_	_	_	‡	‡	‡	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs (the previous value in Rd was H'7F); otherwise cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction increments an 8-bit general register and places the result in the 8-bit general register.

Addressing	Mnem. Operands			Instruction code						
mode	winem.	Operands	1st b	yte	2nd	byte	3rd byte	4th byte	states	
Register direct	INC	Rd	0	Ϊ́Α	0	rd			2	

 $(EAd) \longrightarrow PC$

<Assembly-Language Format>

JMP <EA>

<Examples>

JMP @R6

JMP @H'2000

JMP @@H'9A

<Operand Size>

_

<Condition Code>

I	Н			N	Z	V	C
	_		_	_	_		

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction branches unconditionally to a specified destination address.

The destination address must be even.

Addressing	Mnem. Operands			Instruction code						
mode	willem.	Operanus	1st b	1st byte 2nd byte 3		3rd byte	4th byte	states		
Register indirect	JMP	@Rn	5	9	0 rn	0			4	
Absolute address	JMP	@aa:16	5	Α	0	0	а	ıbs.	6	
Memory indirect	JMP	@@aa:8	5	В	ab	S.			8	

 $PC \rightarrow @-SP$

 $(EAd) \longrightarrow PC$

<Condition Code>

I	Н	N	Z	V	C

<Assembly-Language Format>

JSR <EA>

<Examples>

JSR @R3

JSR @H'1D26

JSR @@H'F0

I: Previous value remains unchanged.

H: Previous value remains unchanged.

N: Previous value remains unchanged.

Z: Previous value remains unchanged.

V: Previous value remains unchanged.

C: Previous value remains unchanged.

<Operand Size>

<Description>

This instruction pushes the program counter onto the stack, then branches to a specified destination address. The program counter value pushed on the stack is the address of the instruction following the JSR instruction. The destination address must be even.

Addressing	Mnem. Operands			Instruction code						
mode	wiriem.	Operanus	1st b	1st byte 2nd byte 3		3rd byte	4th byte	states		
Register indirect	JSR	@Rn	5	D	0 rn	0			6	
Absolute address	JSR	@aa:16	5	E	0	0	a	abs.	8	
Memory indirect	JSR	@@aa:8	5	F	abs.				8	

 $(EAs) \longrightarrow CCR$

<Assembly-Language Format>

LDC <EAs>, CCR

<Examples>

LDC #H'80, CCR

LDC R4H, CCR

<Operand Size>

Byte

<Condition Code>

I		Н			Z	V	C
\$	\leftrightarrow				\$		

- I: Loaded from the source operand.
- H: Loaded from the source operand.
- N: Loaded from the source operand.
- Z: Loaded from the source operand.
- V: Loaded from the source operand.
- C: Loaded from the source operand.

<Description>

This instruction loads the source operand contents into the condition code register (CCR). The source operand can be an 8-bit general register or 8-bit immediate data. Bits 4 and 6 are loaded as well as the flag bits.

No interrupt requests are accepted immediately after this instruction. All interrupts, including the nonmaskable interrupt (NMI), are deferred until after the next instruction.

Addressing	Mnom	Operands		No. of						
mode	Mnem.		1st b	yte	2nd b	yte	3rd byte	4th byte	states	
Immediate	LDC	#xx:8, CCR	0	7	IMM				2	
Register direct	LDC	Rs, CCR	0	3	0	rs			2	

 $Rs \longrightarrow Rd$

<Assembly-Language Format>

MOV.B Rs, Rd

<Examples>

MOV.B R1L, R2H

<Operand Size>

Byte

<Condition Code>

I		Н		N	Z	V	C
_	_	_	_	‡	\$	0	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0."
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction moves one byte of data from a source register to a destination register and sets condition code flags according to the data value.

Addressing	Mnem.	Operanda		No. of			
mode	ivinem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	MOV.B	Rs, Rd	0 C	rs rd			2

 $Rs \longrightarrow Rd$

<Assembly-Language Format>

MOV.W Rs, Rd

<Examples>

MOV.W R3, R4

<Operand Size>

Word

<Condition Code>

	I		Η		N	Z	V	C
-		_	_	_	‡	\$	0	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0."
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction moves one word of data from a source register to a destination register and sets condition code flags according to the data value.

Addressing	Mnem.	Operands	Ins	No. of			
mode	winem.	Operanus	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	MOV.W	Rs, Rd	0 D	0 rs 0 rd			2

 $(EAs) \longrightarrow Rd$

<Assembly-Language Format>

MOV.B <EAs>, Rd

<Examples>

MOV.B @R1, R2H

MOV.B @R5+, R0L

MOV.B @H'FFF1, R1H

MOV.B #H'A5, R3L

<Operand Size>

Byte

<Condition Code>

I		Н	N	Z	V	C
	_			\$	0	

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0"
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction moves one byte of data from a source operand to a destination register and sets condition code flags according to the data value. The source operand can be memory contents or immediate data.

The MOV.B @R7+, Rd instruction should never be used, because it leaves an odd value in the stack pointer. This may result in loss of data, since the stack is always accessed a word at a time at an even address.

Addressing	Mnem.	Operands			Instruct	ion code		No. of	
mode	wiriem.	Operands	1st byte		2nd byte	3rd byte	4th byte	states	
Immediate	MOV.B	#xx:8, Rd	F	rd	IMM			2	
Register indirect	MOV.B	@RS, Rd	6	8	0 rs rd			4	
Register indirect with displacement	MOV.B	@(d:16,Rs),Rd	6	E	0 rs rd		disp.	6	
Register indirect with post-increment	MOV.B	@Rs+, Rd	6	С	0 rs rd			6	
Absolute address	MOV.B	@aa:8, Rd	2	rd	abs			4	
Absolute address	MOV.B	@aa:16, Rd	6	Α	0 rd		abs.	6	

 $(EAs) \rightarrow Rd$

<Assembly-Language Format>

MOV.W <EAs>, Rd

<Examples>

MOV.W @R3, R4

MOV.W @(H'0004,R5), R6

MOV.W @R7+, R0

MOV.W #H'B00A, R1

<Operand Size>

Word

<Condition Code>

I		Н		N	Z	V	C
	_	_	_	‡	\$	0	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0."
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction moves one word of data from a source operand to a destination register and sets condition code flags according to the data value.

If the source operand is in memory, it must be located at an even address.

MOV.W @R7+, Rd is identical in machine language to POP.W Rd.

Addressing	Mnem.	Operands			Instruction	on code		No. of
mode	Willem.	Ореганиз	1st by	yte	2nd byte	3rd byte	4th byte	states
Immediate	MOV.W	#xx:16, Rd	7	9	0 0 rd	[]	MM	4
Register indirect	MOV.W	@RS, Rd	6	9	0 rs 0 rd			4
Register indirect with displacement	MOV.W	@(d:16,Rs),Rd	6	F	0 rs 0 rd		disp.	6
Register indirect with post-increment	MOV.W	@Rs+, Rd	6	D	0 rs 0 rd			6
Absolute address	MOV.W	@aa:16, Rd	6	В	0 0 rd		abs.	6

 $Rs \rightarrow (EAd)$

<Assembly-Language Format>

MOV.B Rs, <EAd>

<Examples>

MOV.B R1L, @R0

MOV.B R3H, @(H'8001, R0)

MOV.B R5H, @-R4

MOV.B R6L, @H'FE77

<Operand Size>

Byte

<Condition Code>

I		Н			Z	V	C
_	_	_		‡		0	

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0."
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction moves one byte of data from a source register to memory and sets condition code flags according to the data value.

The MOV.B Rs, @-R7 instruction should never be used, because it leaves an odd value in the stack pointer. This may result in loss of data, since the stack is always accessed a word at a time at an even address.

The instruction MOV.B RnH, @-Rn or MOV.B RnL, @-Rn decrements register Rn, then moves the upper or lower byte of the decremented result to memory.

Addressing	Mnem. Operands				No. of					
mode	Willem.	Ореганиз	1st byte		2nd byte		oyte	3rd byte	4th byte	states
Register indirect	MOV.B	Rs, @Rd	6	8	1	rd	rs			4
Register indirect with displacement	MOV.B	Rs, @(d:16,Rd)	6	: E	1	rd	rs		disp.	6
Register indirect with pre-decrement	MOV.B	Rs, @-Rd	6	С	1	rs	rs			6
Absolute address	MOV.B	Rs,@aa:8	3	rs		ab	s			4
Absolute address	MOV.B	Rs,@aa:16	6	Α	1	3	rs		abs.	6

 $Rs \rightarrow (EAd)$

<Assembly-Language Format>

MOV.W Rs, <EAd>

<Examples>

MOV.W R3, @R4

MOV.W R2, @(H,0030,R5)

MOV.W R1, @-R7

MOV.W RO, @H'FED6

<Operand Size>

Word

<Condition Code>

I		Н			Z	V	C
_	_			‡	‡	0	

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0."
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction moves one word of data from a general register to memory and sets condition code flags according to the data value.

The destination address in memory must be even.

MOV.W Rs, @-R7 is identical in machine language to PUSH.W Rs.

The instruction MOV.W Rn, @-Rn decrements register Rn by 2, then moves the decremented result to memory.

Addressing	Mnem.	Operands			No. of				
mode	IVIIICIII.	Operanus	1st b	yte	2nd byte	3rd byte 4th byte		states	
Register indirect	MOV.W	Rs, @Rd	6	9	1 rd 0 rs			4	
Register indirect with displacement	MOV.W	Rs, @(d:16, Rd)	6	F	1 rd 0 rs	di	isp.	6	
Register indirect with pre-decrement	MOV.W	Rs, @-Rd	6	D	1 rd 0 rs			6	
Absolute address	MOV.W	Rs, @aa:16	6	В	8 0 rs	al	bs.	6	

synchronization with the E clock (EAs) \rightarrow Rd

<Assembly-Language Format>

MOVFPE @aa:16, Rd

<Examples>

MOVFPE @H'FF81, ROH

<Operand Size>

Byte

<Condition Code>

I		Η	N	Z	V	C
	_			‡	0	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0."
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged

<Description>

This instruction moves one byte of data from an absolute address location to a destination register, and sets the condition code flags according to the data value. The transfer is performed in synchronization with the E (enable) clock used by peripheral devices. The transfer requires 9 to 16 states, so the execution time is variable. For further information on basic timing, See the each *Hardware Manuals*.

This instruction should not be used with chips not having an E clock output pin or in singlechip mode.

When the source operand is located in on-chip memory or the on-chip register field, the MOVFPE instruction is identical in operation to MOV.B @aa:16, Rd.

Note that only 16-bit absolute addressing can be used, and word data cannot be transferred.

Addressing	Mnem.	Operands			No. of				
mode	William.	Ореганаз	1st byte		2nd byte		3rd byte	4th byte	states
Absolute address	MOVFPE	@aa:16, Rd	6	Α	4	rd		abs.	13-20

synchronization with the E clock Rs __ (EAd)

<Assembly-Language Format>

MOVTPE Rs, @aa:16

<Examples>

MOVTPE R2L, @H'FF8D

<Operand Size>

Byte

<Condition Code>

I		Н	N	Z	V	C
_	_		‡	‡	0	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0."
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction moves one byte of data from a source register to an absolute address location, and sets the condition code flags according to the data value. The transfer is performed in synchronization with the E (enable) clock used by peripheral devices. The transfer requires 9 to 16 states, so the execution time is variable. For further information on basic timing, see the each *Hardware Manuals*.

This instruction should not be used with chips not having an E clock output pin or in singlechip mode.

When the destination operand is located in on-chip memory or the on-chip register field, the MOVTPE instruction is identical in operation to MOV.B Rs, @aa:16.

Note that only 16-bit absolute addressing can be used, and word data cannot be transferred.

Addressing	Mnem. Operands				No. of		
mode	winem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Absolute address	MOVTPE	Rs, @aa:16	6 A	C rs	8	abs.	13-20

 $Rd \times Rs \rightarrow Rd$

<Assembly-Language Format>

MULXU Rs, Rd

<Examples>

MULXU ROH, R3

<Operand Size>

Byte

<Condition Code>

I		Н	N	Z	V	C
_	_			_		_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction performs 8-bit \times 8-bit \to 16-bit multiplication. It multiplies a destination register by a source register and places the result in the destination register. The source register is an 8-bit register. The destination register is a 16-bit register containing the data to be multiplied in the lower byte. (The upper byte is ignored). The result is placed in both bytes of the destination register. The operation is shown schematically below.

The multiplier can occupy either the upper or lower byte of the source register.

Addressing	Mnem.	Operands	Inst	No. of			
mode	iviriem.	Operanus	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	MULXU	Rs, Rd	5 0	rs 0 rd			14

 $0 - Rd \longrightarrow Rd$

<Assembly-Language Format>

NEG Rd

<Examples>

NEG ROL

<Operand Size>

Byte

<Condition Code>

	I	Н		N	Z	V	C
ſ		\$	_		\$	\$	\$

- I: Previous value remains unchanged.
- H: Set to "1" when there is a borrow from bit 3; otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs (the previous contents of the destination register was H'80); otherwise cleared to "0."
- C: Set to "1" if there is a borrow from bit 7 (the previous contents of the destination register was not H'00); otherwise cleared to "0."

<Description>

This instruction replaces the contents of an 8-bit general register with its two's complement. (subtracts the register contents from H'00).

If the original contents of the destination register was H'80, the register value remains H'80 and the overflow flag is set.

Addressing mode	Mnem.	Operands	Inst	No. of			
			1st byte	2nd byte	3rd byte	4th byte	states
Register direct	NEG	Rd	1 7	8 rd			2

 $PC + 2 \rightarrow PC$

<Assembly-Language Format>

NOP

<Examples>

NOP

<Operand Size>

<Condition Code>

I	Η		N	Z	V	C
_		_	_	_	_	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction only increments the program counter, causing the next instruction to be executed. The internal state of the CPU does not change.

The NOP instruction can be used to fill in gaps in programs, or for software synchronization.

Addressing	Mnem. Operands			No. of						
mode	winem.	Operands	1st b	yte	2nd b	oyte	3rd byte	4th byte	states	
	NOP		0	0	0	0			2	

 $\neg Rd \xrightarrow{} Rd$

<Assembly-Language Format>

NOT Rd

<Examples>

NOT R4L

<Operand Size>

Byte

<Condition Code>

I	Н		N	Z	V	C
_	 _	_	‡	\$	0	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction replaces the contents of an 8-bit general register with its one's complement (subtracts the register contents from H'FF).

Addressing	Mnem. Operands		Inst		No. of		
mode	winem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	NOT	Rd	1 7	0 rd			2

 $Rd_{\vee}(EAs) \longrightarrow Rd$

<Assembly-Language Format>

OR <EAs>, Rd

<Examples>

OR R2H, R3H

OR #H'CO, ROH

<Operand Size>

Byte

<Condition Code>

I		Н	N	Z	V	C
	_	_		\$	0	

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction ORs the source operand with the contents of an 8-bit general register and places the result in the general register.

The source operand can be an 8-bit register value or immediate byte data.

Addressing	Mnem.	Operands		Instruction code					
mode	iviriem.	Operands	1st byte		2nd byte		3rd byte	4th byte	states
Immediate	OR	#xx:8, Rd	С	rd	IIV	1M			2
Register direct	OR	Rs, Rd	1	4	rs	rd			2

 $CCR \downarrow \#IMM \rightarrow CCR$

<Assembly-Language Format>

ORC #xx:8, CCR

<Examples>

ORC #H'80, CCR

<Operand Size>

Byte

<Condition Code>

I		Н		N	Z	V	C
\leftrightarrow					\$		\leftrightarrow

I: ORed with bit 7 of the immediate data.

H: ORed with bit 5 of the immediate data.

N: ORed with bit 3 of the immediate data.

Z: ORed with bit 2 of the immediate data.

V: ORed with bit 1 of the immediate data.

C: ORed with bit 0 of the immediate data.

<Description>

This instruction ORs the condition code register (CCR) with immediate data and places the result in the condition code register. Bits 6 and 4 are ORed as well as the flag bits.

No interrupt requests are accepted immediately after this instruction. All interrupts, including the nonmaskable interrupt (NMI), are deferred until after the next instruction.

Addressing mode	Mnem.	Operande		No. of				
	winem.	Operands	1st byte		2nd byte	3rd byte	4th byte	states
Immediate	ORC	#xx:8, CCR	0	4	IMM			2

 $@SP+ \rightarrow Rn$

<Assembly-Language Format>

POP Rn

<Examples>

POP R1

<Operand Size>

Word

<Condition Code>

I		Н		N	Z	V	C
	_	_	_	‡	‡	0	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0."
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction pops data from the stack to a 16-bit general register and sets condition code flags according to the data value.

POP.W Rn is identical in machine language to MOV.W @SP+, Rn.

Addressing	ssing Mnem. Op				I		No. of			
mode	Milletti.	Operands	1st b	yte	2nd	byte	3rd byte	4th byte	states	
	POP	Rd	6	D	7	0 rn			6	

Rn _→ @–SP

<Assembly-Language Format>

PUSH Rn

<Examples>

PUSH R2

<Operand Size>

Word

<Condition Code>

I	Η	N	Z	V	C
_	 _	 ‡	‡	0	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the data value is negative; otherwise cleared to "0."
- Z: Set to "1" if the data value is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction pushes data from a 16-bit general register onto the stack and sets condition code flags according to the data value.

PUSH.W Rn is identical in machine language to MOV.W Rn, @-SP.

Addressing mode	Mnom	Mnem. Operands		Instruction code					
	winem.	Operands	1st byte	2nd byte	3rd byte	4th byte	No. of states		
	PUSH	Rs	6 D	F 0 rn			6		

ROTL (ROTate Left)

ROTL

<Operation>

Rd (rotated left) _ Rd

<Assembly-Language Format>

ROTL Rd

<Examples>

ROTL R2L

<Operand Size>

Byte

<Condition Code>

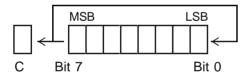
I		Н	N	Z	V	C
_	_			\$	0	\leftrightarrow

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Receives the previous value in bit 7.

<Description>

This instruction rotates an 8-bit general register one bit to the left. The most significant bit is rotated to the least significant bit, and also copied to the carry flag.

The operation is shown schematically below.



Addressing mode	Mnem.	Operands			No. of		
	winem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	ROTL	Rd	1 2	8 rd			2

Rd (rotated right) _ Rd

<Assembly-Language Format>

ROTR Rd

<Examples>

ROTR R5L

<Operand Size>

Byte

<Condition Code>

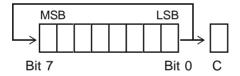
Ι	Η		N	Z	V	C
_		_	\Rightarrow	\$	0	1

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Receives the previous value in bit 0.

<Description>

This instruction rotates an 8-bit general register one bit to the right. The least significant bit is rotated to the most significant bit, and also copied to the carry flag.

The operation is shown schematically below.



Addressing mode	Mnem.	Operands		No. of			
			1st byte	2nd byte	3rd byte	4th byte	states
Register direct	ROTR	Rd	1 3	8 rd			2

Rd (rotated with carry left) \rightarrow Rd

<Assembly-Language Format>

ROTXL Rd

<Examples>

ROTXL R1H

<Operand Size>

Byte

<Condition Code>

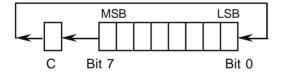
I		Η	N	Z	V	C
_	_	_	 	‡	0	‡

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Receives the previous value in bit 7.

<Description>

This instruction rotates an 8-bit general register one bit to the left through the carry flag. The carry flag is rotated into the least significant bit of the register. The most significant bit rotates into the carry flag.

The operation is shown schematically below.



Addressing mode	Mnem.	Operands		No. of			
			1st byte	2nd byte	3rd byte	4th byte	states
Register direct	ROTXL	Rd	1 2	0 rd			2

Rd (rotated with carry right) \rightarrow Rd

<Assembly-Language Format>

ROTXR Rd

<Examples>

ROTXR R5L

<Operand Size>

Byte

<Condition Code>

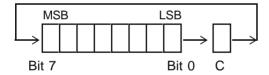
I		Н		N	Z	V	C
	_		_		\$	0	

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Receives the previous value in bit 0.

<Description>

This instruction rotates an 8-bit general register one bit to the right through the carry flag. The least significant bit is rotated into the carry flag. The carry flag rotates into the most significant bit.

The operation is shown schematically below



Addressing	Mnem.	Operanda		Instruction	on code		No. of
mode	iviriem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	ROTXR	Rd	1 3	0 rd			2

 $@SP+ \rightarrow CCR$

 $@SP+ \longrightarrow PC$

<Condition Code>

<Assembly-Language Format>

RTE

<Examples>

RTE

<Operand Size>

I: Restored from stack.

H: Restored from stack.

N: Restored from stack.

Z: Restored from stack.

V: Restored from stack.

C: Restored from stack.

<Description>

This instruction returns from an interrupt-handling routine. It pops the condition code register (CCR) and program counter (PC) from the stack. Program execution continues from the address restored to the program counter.

The CCR and PC contents at the time of execution of this instruction are lost.

The CCR is one byte in size, but it is popped from the stack as a word (in which the lower 8 bits are ignored).

This instruction therefore adds 4 to the value of the stack pointer (R7).

Addressing	Mnem.	Operands		Instructi	on code		No. of
mode	iviriem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
	RTE		5 6	7 0			10

 $@SP+ \longrightarrow PC$

<Assembly-Language Format>

RTS

<Examples>

RTS

<Operand Size>

<Condition Code>

I		Η	N	Z	V	C
_	_		 _			

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Previous value remains unchanged.
- Z: Previous value remains unchanged.
- V: Previous value remains unchanged.
- C: Previous value remains unchanged.

<Description>

This instruction returns from a subroutine. It pops the program counter (PC) from the stack.

Program execution continues from the address restored to the program counter.

The PC contents at the time of execution of this instruction are lost.

Addressing	Mnem.	Operands		Instructi	on code		No. of
mode	iviriem.	Operands	1st byte	2nd byte	2nd byte 3rd byte 4th by		states
	RTS		5 4	7 0			8

Rd (shifted arithmetic left) → Rd

<Assembly-Language Format>

SHAL Rd

<Examples>

SHAL R5H

<Operand Size>

Byte

<Condition Code>

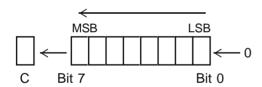
_	I	Η	N	Z	V	C
ĺ				\$	\$	\$

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Receives the previous value in bit 7.

<Description>

This instruction shifts an 8-bit general register one bit to the left. The most significant bit shifts into the carry flag, and the least significant bit is cleared to "0."

The operation is shown schematically below.



The SHAL instruction is identical to the SHLL instruction except for its effect on the overflow (V) flag.

Addressing mode	Mnom	Operanda		Instructi	on code		No. of
	Mnem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	SHAL	Rd	1 0	8 rd			2

Rd (shifted arithmetic right) → Rd

<Assembly-Language Format>

SHAR Rd

<Examples>

SHAR R5H

<Operand Size>

Byte

<Condition Code>

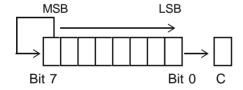
I		Η	N	Z	V	C
_	_				0	

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Receives the previous value in bit 0.

<Description>

This instruction shifts an 8-bit general register one bit to the right. The most significant bit remains unchanged. The sign of the result does not change. The least significant bit shifts into the carry flag.

The operation is shown schematically below.



Addressing mode	Mnom	Operanda		Instructi	on code		No. of
	Mnem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	SHAR	Rd	1 1	8 rd			2

Rd (shifted logical left) \rightarrow Rd

<Assembly-Language Format>

SHLL Rd

<Examples>

SHLL R2L

<Operand Size>

Byte

<Condition Code>

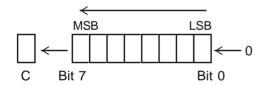
I	Η	N	Z	V	C
	 		‡	0	\$

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Receives the previous value in bit 0.

<Description>

This instruction shifts an 8-bit general register one bit to the left. The least significant bit is cleared to "0." The most significant bit shifts into the carry flag.

The operation is shown schematically below.



The SHLL instruction is identical to the SHAL instruction except for its effect on the overflow (V) flag.

Addressing	Mnem. Operands			No. of				
mode	willetti.	Operanus	1st byte	2nd byte	3rd byte	4th byte	states	
Register direct	SHLL	Rd	1 0	0 rd			2	

Rd (shifted logical right) \rightarrow Rd

<Assembly-Language Format>

SHLR Rd

<Examples>

SHLR R3L

<Operand Size>

Byte

<Condition Code>

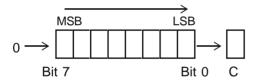
I		Н	N	Z	V	C
	_			‡	0	\leftrightarrow

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" if the result is negative; otherwise cleared to "0."
- Z: Set to "1" if the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Receives the previous value in bit 0.

<Description>

This instruction shifts an 8-bit general register one bit to the right. The most significant bit is cleared to 0. The least significant bit shifts into the carry flag.

The operation is shown schematically below.



Addressing mode	Mnom	Mnem. Operands		Instruction code				
	winem.	Operands	1st byte	2nd byte	3rd byte	4th byte	No. of states	
Register direct	SHLR	Rd	1 1	0 rd			2	

 Condition Code

 Program execution state → power-down mode
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C
 I H N Z V C

 I H N Z V C C
 I H N Z V C

 I H Previous value remains unchanged.

 I H Previous value remai

SLEEP

<Description>

SLEEP (SLEEP)

When the SLEEP instruction is executed, the CPU enters a power-down mode. Its internal state remains unchanged, but the CPU stops executing instructions and waits for an exception-handling request (interrupt or reset). When it receives an exception-handling request, the CPU exits the power-down mode and begins the exception-handling sequence.

If the interrupt mask (I) bit is set to "1," the power-down mode can be released only by a nonmaskable interrupt (NMI) or reset.

For information about the power-down modes, see the *Hardware Manual* for the particular chip.

Addressing mode	Mnom	Mnem. Operands		Instruction code				
	willem.	Operanus	1st byte	2nd byte	3rd byte	4th byte	states	
	SLEEP		0 1	8 0			2	

 $CCR \rightarrow Rd$

<Assembly-Language Format>

STC CCR, Rd

<Examples>

STC CCR, R6H

<Operand Size>

Byte

<Condition Code>

I	Η	N	Z	V	C
					_

I: Previous value remains unchanged.

H: Previous value remains unchanged.

N: Previous value remains unchanged.

Z: Previous value remains unchanged.

V: Previous value remains unchanged.

C: Previous value remains unchanged.

<Description>

This instruction copies the condition code register (CCR) to a specified general register. Bits 6 and 4 are copied as well as the flag bits.

Addressing mode	Mnom	Operanda		No. of			
	Mnem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	STC	CCR, Rd	0 2	0 rd			2

 $Rd - Rs \longrightarrow Rd$

<Assembly-Language Format>

SUB.B Rs, Rd

<Examples>

SUB.B ROL, R2L

<Operand Size>

Byte

<Condition Code>

I		Η	N	Z	V	C
_	_		\leftrightarrow	\$		\$

- I: Previous value remains unchanged.
- H: Set to "1" when there is a borrow from bit 3; otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Set to "1" if there is a borrow from bit 7; otherwise cleared to "0."

<Description>

This instruction subtracts an 8-bit source register from an 8-bit destination register and places the result in the destination register.

Only register direct addressing is supported. To subtract immediate data it is necessary to use the SUBX.B instruction, first setting the zero flag to "1" and clearing the carry flag to "0".

The following codings can also be used to subtract nonzero immediate data.

(1) ORC #H'05, CCR SUBX #(Imm-1), Rd (2) ADD #(0-Imm), Rd XORC #H'01, CCR

Addressing	Mnom	Mnem. Operands		Instruction code				
mode	iviriem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states	
Register direct	SUB.B	Rs, Rd	1 8	rs rd			2	

 $Rd - Rs \rightarrow Rd$

<Assembly-Language Format>

SUB.W Rs, Rd

<Examples>

SUB.W RO, R1

<Operand Size>

Word

<Condition Code>

I		Η	N	Z	V	C
	_		\Rightarrow			\leftrightarrow

- I: Previous value remains unchanged.
- H: Set to "1" when there is a borrow from bit 11; otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Set to "1" if there is a borrow from bit 15; otherwise cleared to "0."

<Description>

This instruction subtracts a 16-bit source register from a 16-bit destination register and places the result in the destination register.

Addressing	Mnem.	Operands	Instructi		on code	No. of	
mode	winem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Register direct	SUB.W	Rs, Rd	1 9	0 rs 0 rd			2

SUBS (SUBtract with Sign extension)

SUBS

<Operation>

 $Rd - 1 \longrightarrow Rd$

 $Rd-2 \underset{\longrightarrow}{\longrightarrow} Rd$

<Condition Code>

<Assembly-Language Format>

SUBS #1, Rd

SUBS #2, Rd

H: Previous value remains unchanged.

N: Previous value remains unchanged.

I: Previous value remains unchanged.

Z: Previous value remains unchanged.

V: Previous value remains unchanged.

C: Previous value remains unchanged.

<Examples>

SUBS #1, R3

SUBS #2, R5

<Operand Size>

Word

<Description>

This instruction subtracts the immediate value 1 or 2 from word data in a general register.

Differing from the SUB instruction, it does not affect the condition code flags.

The SUBS instruction does not permit byte operands.

Addressing	Mnem. Operands -			Instruction code				
mode	WITIETTI.	Operands	1st byte	2nd byte	3rd byte	4th byte	states	
Register direct	SUBS	#1, Rd	1 B	0 0 rd			2	
Register direct	SUBS	#2, Rd	1 B	8 0 rd			2	

 $Rd - (EAs) - C \longrightarrow Rd$

<Assembly-Language Format>

SUBX <EAs>, Rd

<Examples>

SUBX ROL, R3L SUBX #H'32, R5H

<Operand Size>

Byte

<Condition Code>

I		Н	N	Z	V	C
_	_	\$	‡	\$		‡

- I: Previous value remains unchanged.
- H: Set to "1" if there is a borrow from bit 3; otherwise cleared to "0."
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Previous value remains unchanged when the result is zero; otherwise cleared to "0"
- V: Set to "1" if an overflow occurs; otherwise cleared to "0."
- C: Set to "1" if there is a borrow from bit 7; otherwise cleared to "0."

<Description>

This instruction subtracts the source operand and carry flag from the contents of an 8-bit general register and places the result in the general register.

The source operand can be an 8-bit register value or immediate byte data.

Addressing	Mnem. Operands			Instruction code				
mode	winem.	Operanus	1st byte	2nd byte	3rd byte	4th byte	states	
Immediate	SUBX	#xx:8, Rd	B rd	IMM			2	
Register direct	SUBX	Rs, Rd	1 E	rs rd			2	

 $Rd \oplus (EAs) \rightarrow Rd$

<Assembly-Language Format>

XOR <EAs>, Rd

<Examples>

XOR ROH, R1H XOR #H'F0, R2L

<Operand Size>

Byte

<Condition Code>

I	Η		N	Z	V	C
	 _	_		\$	0	_

- I: Previous value remains unchanged.
- H: Previous value remains unchanged.
- N: Set to "1" when the result is negative; otherwise cleared to "0."
- Z: Set to "1" when the result is zero; otherwise cleared to "0."
- V: Cleared to "0."
- C: Previous value remains unchanged.

<Description>

This instruction exclusive-ORs the source operand with the contents of an 8-bit general register and places the result in the general register.

The source operand can be an 8-bit register value or immediate byte data.

Addressing	Mnom	Operanda		Instructi	on code		No. of
mode	Mnem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Immediate	XOR	#xx:8, Rd	D rd	IMM			2
Register direct	XOR	Rs, Rd	1 5	rs rd			2

 $CCR \oplus \#IMM \longrightarrow CCR$

<Assembly-Language Format>

XORC #xx:8, CCR

<Examples>

XORC #H'50, CCR

<Operand Size>

Byte

<Condition Code>

I		Η		N	Z	V	C
\Rightarrow	↑	\$		1	\$	\updownarrow	\leftrightarrow

- I: Exclusive-ORed with bit 7 of the immediate data.
- H: Exclusive-ORed with bit 5 of the immediate data.
- N: Exclusive-ORed with bit 3 of the immediate data.
- Z: Exclusive-ORed with bit 2 of the immediate data.
- V: Exclusive-ORed with bit 1 of the immediate data.
- C: Exclusive-ORed with bit 0 of the immediate data.

<Description>

This instruction exclusive-ORs the condition code register (CCR) with immediate data and places the result in the condition code register. Bits 6 and 4 are exclusive-ORed as well as the flag bits.

No interrupt requests are accepted immediately after this instruction. All interrupts, including the nonmaskable interrupt (NMI), are deferred until after the next instruction.

Addressing	Mnem.	Operands		Instructi	on code		No. of
mode	iviriem.	Operands	1st byte	2nd byte	3rd byte	4th byte	states
Immediate	XORC	#xx:8, CCR	0 5	IMM			2

Appendix A. Operation Code Map

The table in this appendix is a map of the operation codes contained in the first byte of the instruction code (bits 15 to 8 of the first instruction word). Some pairs of instructions have indentical first bytes. These instructions are differentiated by the first bit of the second byte (bit 7 of the first instruction word).

Instruction when first bit of byte 2 (bit 7 of first instruction word) is "0." Instruction when first bit of byte 2 (bit 7 of first instruction word) is "1."

Operation Code Map

H/O	0	1 S	2	3	4	5	9	7	8	6	A	В	C	D	田	Ī
0	NOP	SHLL			BRA	MULXU	TING	bae1								
1	SLEEP	SHLR			BRN	MULXU DIVXU	TOMO	DINOI								
2	STC	ROTXL			BHI		_	BCLK								
3	LDC	ROTXR			BLS			B151								
4	ORC	OR			BCC	RTS		BORBIOR								
5	XORC	XOR			BCS	BSR		BXOR BIXOR								
9	ANDC	AND			BNE	RTE		BAND BIAND								
7	LDC	NOT		MO	BEQ		BST BIST		ADD	ADDX	CMP	SUBX	OR	XOR	AND	MOV
8	ADD	SUB			BVC											
6	Q	B			BVS			MOV								
А	INC	DEC			BPL	JMP										
В	ADDS	SUBS			BMI		W	EEPMOV								
C	M	C			BGE		MOV*									
О	MOV	CMP			BLT			nanipulat								
Е	ADDX	SUBX			BGT	JSR		Bit manipulation instruction								
Н	DAA	DAS			BLE			ction								

^{*} The MOVFPE and MOVTPE instructions are identical to MOV instructions in the first byte and first bit of the second byte (bits 15 to 7 of the instruction word). They are differentiated in the next bit (bit 6 of the instruction word). See section 2.

Appendix. B Instruction Set List

						lres												
Mnemonic	Operand size	Operation	#xx: 8/16		u	@(d:16, Rn)	@-Rn/@Rn+	@aa: 8/16	@(d:8, PC)	aa	Implied	(Con	diti	on	cod	e	of states
	Ор		:xx#	Rn	@Rn	p)@	<u>@</u> -1	@ a s	p)@	@ @ aa	Imp	I	Н	N	Z	V	C	No.
MOV.B #xx:8,Rd	В	#xx:8→Rd8	2									_	-	1	1	0	_	2
MOV.B Rs,Rd	В	$Rs8 \rightarrow Rd8$		2								_	-	1	1	0	_	2
MOV.B @Rs,Rd	В	@Rs16→Rd8			2							_	_	1	1	0	_	4
MOV.B @(d:16,Rs),Rd	В	$@(d:16,Rs16) \rightarrow Rd8$				4						_	_	1	1	0	-	6
MOV.B @Rs+,Rd	В	@Rs16→Rd8 Rs16+1→Rs16					2					-	_	‡	1	0	-	6
MOV.B @aa:8,Rd	В	@aa:8→Rd8						2				_	-	1	1	0	_	4
MOV.B @aa:16,Rd	В	@aa:16→Rd8						4				_	_	1	1	0	_	6
MOV.B Rs,@Rd	В	Rs8→@Rd16			2							_	-	1	1	0	-	4
MOV.B Rs,@(d:16,Rd)	В	Rs8→@(d:16,Rd16)				4						-	-	1	1	0	-	6
MOV.B Rs,@-Rd	В	Rd16–1 → Rd16 Rs8 → @Rd16					2					_	_	‡	‡	0	_	6
MOV.B Rs,@aa:8	В	Rs8→@aa:8						2				_	_	1	1	0	_	4
MOV.B Rs,@aa:16	В	Rs8→@aa:16						4				_	_	1	1	0	_	6
MOV.W #xx:16,Rd	w	#xx:16 → Rd	4									_	_	Ì	1	0	_	4
MOV.W Rs,Rd	w	Rs16→Rd16		2								_	_	1	1	0	_	2
MOV.W @Rs,Rd	w	@Rs16→Rd16			2							_	_	1	1	0	_	4
MOV.W @(d:16,Rs),Rd	W	@(d:16,Rs16) →Rd16				4						_	_	1	1	0	_	6
MOV.W @Rs+,Rd	w	@Rs16→Rd16					2					_	_	†	†	0	_	6
,		Rs16+2→Rs16																
MOV.W @aa:16,Rd	w	@aa:16→Rd16						4				_	-	1	1	0	_	6
MOV.W Rs,@Rd	w	Rs16→@Rd16			2							_	_	1	1	0	_	4
MOV.W Rs,@(d:16,Rd)	w	Rs16→@(d:16,Rd16)				4						_	_	1	1	0	-	6
MOV.W Rs,@-Rd	W	Rs16-2→Rd16 Rs16→@Rd16					2					-	-	‡	‡	0	_	6
MOV.W Rs,@aa:16	w	Rs16→@aa:16						4				_	_	1	1	0	_	6
POP Rd	W	@SP→Rd16 SP+2→SP					2					-	-	†	†	0	_	6
PUSH Rs	W	$\begin{array}{c} SP-2 \rightarrow SP \\ Rs16 \rightarrow @SP \end{array}$					2					-	-	‡	1	0	-	6
MOVFPE @aa:16,Rd	В	@aa:16→Rd (Synchronization with E clock)						4				_	_	‡	1	0	_	(5)
MOVTPE Rs,@aa:16	В	Rs→@aa:16 (Synchronization with E clock)						4				-	-	1	1	0	-	(5)

						essi ıcti											
Mnemonic	Operand size	Operation	8/16		u	@(d:16, Rn)	@-Rn/@Rn+	@aa: 8/16	@(d:8, PC)	aa		Con	diti	on (cod	e	of states
	obo		#xx:	Rn	@ Rn	p)@	<u>@</u>	@ a a	p)@	@ @	I	Н	N	Z	V	C	No. 0
ADD.B #xx:8,Rd	В	Rd8+#xx:8 → Rd8	2								_	‡	1	1	‡	1	2
ADD.B Rs,Rd	В	Rs8+Rd8 →Rd8		2							-	1	‡	1	‡	1	2
ADD.W Rs,Rd	W	Rs16+Rd16 → Rd16		2							_	1	1	1	1	1	2
ADDX.B #xx:8,Rd	В	Rd8+#xx:8+C → Rd8	2								_	1	1	2	‡	1	2
ADDX.B Rs,Rd	В	Rd8+Rs8+C →Rd8		2							-	1	1	2	‡	1	2
ADDS.W #1,Rd	W	Rd16+1 → Rd16		2							-	_	_	-	ı	_	2
ADDS.W #2,Rd	W	Rd16+2→Rd16		2							-	-	-	_	١	_	2
INC.B Rd	В	Rd8+1→Rd8		2							-	-	1	1	1	_	2
DAA.B Rd	В	Rd8 decimal adjust → Rd8		2							-	*	1	1	*	3	2
SUB.B Rs,Rd	В	Rd8-Rs8 → Rd8		2							-	1	1	1	1	1	2
SUB.W Rs,Rd	W	Rd16-Rs16 → Rd16		2							_	1	1	1	1	1	2
SUBX.B #xx:8,Rd	В	Rd8-#xx:8-C → Rd8	2								-	1	1	2	1	1	2
SUBX.B Rs,Rd	В	Rd8-Rs8-C →Rd8		2							-	1	1	2	‡	1	2
SUBS.W #1,Rd	W	Rd16-1 → Rd16		2							-	-	-	-	_	-	2
SUBS.W #2,Rd	W	Rd16-2 → Rd16		2							-	-	_	-	_	-	2
DEC.B Rd	В	Rd8–1 → Rd8		2							-	_	1	1	1	-	2
DAS.B Rd	В	Rd8 decimal adjust → Rd8		2							_	*	1	1	*	_	2
NEG.B Rd	В	0-Rd →Rd		2							-	1	1	1	1	1	2
CMP.B #xx:8,Rd	В	Rd8-#xx:8	2								_	‡	‡	1	1	‡	2
CMP.B Rs,Rd	В	Rd8-Rs8		2							-	1	1	1	1	‡	2
CMP.W Rs,Rd	W	Rd16-Rs16		2							_	1	1	1	1	1	2
MULXU.B Rs,Rd	В	Rd8×Rs8→Rd16		2							_	-	_	_	-	_	14
DIVXU.B Rs,Rd	В	Rd16+Rs8 → Rd16		2							-	-	6	7	_	-	14
		(RdH:remainder,RdL:quotient)															
AND.B #xx:8,Rd	В	Rd8∧#xx:8→Rd8	2								-	-	1	1	0	_	2
AND.B Rs,Rd	В	$Rd8 \land Rs8 \rightarrow Rd8$		2							-	_	1	1	0	-	2
OR.B #xx:8,Rd	В	Rd8∨#xx:8 → Rd8	2								-	_	1	1	0	_	2
OR.B Rs,Rd	В	Rd8∨Rs8 → Rd8		2							-	-	1	1	0	_	2
XOR.B #xx:8,Rd	В	Rd8⊕#xx:8→Rd8	2								-	-	1	1	0	_	2
XOR.B Rs,Rd	В	Rd8⊕Rs8 → Rd8		2							-	-	1	1	0	_	2
NOT.B Rd	В	Rd →Rd		2							-	-	1	1	0	_	2

				A	ddı ıstr	ess	ing on	mo len	de/ gth								
Mnemonic	Operand size	Operation	8/16			@(d:16, Rn)	@-Rn/@Rn+	@aa: 8/16	@(d:8, PC)	ra		Con	diti	on	cod	e	No. of states
	Ope		#xx: 8/16	Rn	@Rn	@(d :	@- R	@aa	@(q :	@ @ aa	I	Н	N	Z	v	C	No. 0
SHAL.B Rd	В	C ← 0 ← 0 bo		2							_	-	1	1	‡	1	2
SHAR.B Rd	В	b ₇ C b ₀		2							_	-	1	1	0	1	2
SHLL.B Rd	В	C ← ← 0 b ₇ b ₀		2							_	-	‡	1	0	1	2
SHLR.B Rd	В	$0 \xrightarrow{b_7} C$		2							-	_	0	1	0	1	2
ROTXL.B Rd	В	D ₇ D ₀ D ₀		2							-	-	‡	1	0	‡	2
ROTXR.B Rd	В	b ₇ b ₀		2							_	_	‡	1	0	‡	2
ROTL.B Rd	В	b ₇ b ₀		2								-	‡	1	0	‡	2
ROTR.B Rd	В	b_7 b_0		2							_	-	1	1	0	1	2
BSET #xx:3,Rd	В	(#xx:3 of Rd8) ←1		2							-	-	-	-	-	-	2
BSET #xx:3,@Rd	В	(#xx:3 of @Rd16) ←1			4						-	-	-	-	-	-	8
BSET #xx:3,@aa:8	В	(#xx:3 of @aa:8) ←1						4			-	-	-	-	-	-	8
BSET Rn,Rd	В	(Rn8 of Rd8) ←1		2							-	-	-	-	-	_	2
BSET Rn,@Rd	В	(Rn8 of @Rd16) ←1			4						_	_	_	_	_	_	8
BSET Rn,@aa:8	В	(Rn8 of @aa:8) ←1						4			_	_	_	_	_	_	8
BCLR #xx:3,Rd	В	(#xx:3 of Rd8) ← 0		2							_	_	_	_	_	_	2
BCLR #xx:3,@Rd	В	(#xx:3 of @Rd16) ← 0			4						_	_	_	_	_	_	8
BCLR #xx:3,@aa:8	В	(#xx:3 of @aa:8) ← 0						4			_	_	_	_	_	_	8
BCLR Rn,Rd	В	(Rn8 of Rd8) ← 0		2							_	_	_	_	_	_	2
BCLR Rn,@Rd	В	(Rn8 of @Rd16) ←0			4						_	_	_	_	_	_	8
BCLR Rn,@aa:8	В	(Rn8 of @aa:8) ← 0						4			_	-	_	_	_	_	8
BNOT #xx:3,Rd	В	(#xx:3 of Rd8) ← (#xx:3 of Rd8)		2							_	_	_	_	_	_	2
BNOT #xx:3,@Rd	В	(#xx:3 of @Rd16) ← (#xx:3 of @Rd16)			4						_	_	_	_	_	_	8
BNOT #xx:3,@aa:8	В	(#xx:3 of @aa:8) ← (#xx:3 of @aa:8)						4			_	_	_	_	_	_	8

				Ao	ddr stri	essi ıcti	ng on l	mo eng	de/ gth								
Mnemonic	Operand size	Operation	8/16		n	@(d:16, Rn)	@- R n/@ R n+	@aa: 8/16	@(d:8, PC)	aa	(Con	diti	on (cod	e	of states
	Ope		#xx:	Rn	@Rn	p)@	@-F	@aa	p)@	@ @	I	Н	N	Z	V	C	No. 0
BNOT Rn,Rd	В	$(Rn8 \text{ of } Rd8) \leftarrow (\overline{Rn8 \text{ of } Rd8})$		2							_	_	-	_	١	-	2
BNOT Rn,@Rd	В	$(Rn8 \text{ of } @Rd16) \leftarrow \overline{(Rn8 \text{ of } @Rd16)}$			4						-	_	-	_	-	-	8
BNOT Rn,@aa:8	В	$(Rn8 \text{ of } @aa:8) \leftarrow \overline{(Rn8 \text{ of } @aa:8)}$						4			_	_	_	-	ı	-	8
BTST #xx:3,Rd	В	$\overline{(\#xx:3 \text{ of } \text{Rd8})} \rightarrow Z$		2							_	_	_	1	-	_	2
BTST #xx:3,@Rd	В	(#xx:3 of @Rd16)→Z			4						_	-	-	1	ı	-	6
BTST #xx:3,@aa:8	В	(#xx:3 of @aa:8) → Z						4			-	_	-	1	_	_	6
BTST Rn,Rd	В	$\overline{(Rn8 \text{ of } Rd8)} \rightarrow Z$		2							-	-	-	1	-	_	2
BTST Rn,@Rd	В	(Rn8 of @Rd16)→Z			4						_	-	_	1	-	_	6
BTST Rn,@aa:8	В	(Rn8 of @aa:8) → Z						4			-	-	-	1	-	-	6
BLD #xx:3,Rd	В	(#xx:3 of Rd8)→C		2							-	-	-	-	-	1	2
BLD #xx:3,@Rd	В	(#xx:3 of @Rd16)→C			4						-	-	-	-	-	‡	6
BLD #xx:3,@aa:8	В	(#xx:3 of @aa:8)→C						4			-	-	-	-	-	‡	6
BILD #xx:3,Rd	В	(#xx:3 of Rd8)→C		2							-	-	-	-	_	1	2
BILD #xx:3,@Rd	В	(#xx:3 of @Rd16)→C			4						-	-	-	_	-	‡	6
BILD #xx:3,@aa:8	В	(#xx:3 of @aa:8) → C						4			_	-	-	_	-	‡	6
BST #xx:3,Rd	В	C→(#xx:3 of Rd8)		2							-	-	-	_	-	_	2
BST #xx:3,@Rd	В	C→(#xx:3 of @Rd16)			4						_	_	_	_	-	_	8
BST #xx:3,@aa:8	В	C→(#xx:3 of @aa:8)						4			-	-	-	_	_	_	8
BIST #xx:3,Rd	В	$\overline{C} \rightarrow (\#xx:3 \text{ of Rd8})$		2							_	-	-	_	_	-	2
BIST #xx:3,@Rd	В	$\overline{C} \rightarrow (\#xx:3 \text{ of } @Rd16)$			4						-	-	-	_	-	-	8
BIST #xx:3,@aa:8	В	<u>C</u> →(#xx:3 of @aa:8)						4			-	-	-	_	_	-	8
BAND #xx:3,Rd	В	$C \land (\#xx:3 \text{ of } Rd8) \rightarrow C$		2							_	-	_	_	_	1	2
BAND #xx:3,@Rd	В	C^(#xx:3 of @Rd16)→C			4						_	-	_	_	_	†	6
BAND #xx:3,@aa:8	В	$C \land (\#xx:3 \text{ of } @aa:8) \rightarrow C$						4			-	-	_	_	_	1	6
BIAND #xx:3,Rd	В	$C \wedge (\overline{\#xx:3 \text{ of } Rd8}) \rightarrow C$		2							_	-	_	_	_	1	2
BIAND #xx:3,@Rd	В	$C \land \overline{(\#xx:3 \text{ of } @ \text{Rd16})} \rightarrow C$			4						-	-	_	_	_	‡	6
BIAND #xx:3,@aa:8	В	$C \wedge \overline{(\#xx:3 \text{ of } @aa:8)} \rightarrow C$						4			_	_	_	_	_	‡	6
BOR #xx:3,Rd	В	$C \lor (\#xx:3 \text{ of } Rd8) \rightarrow C$		2							_	-	_	_	_	1	2
BOR #xx:3,@Rd	В	C∨(#xx:3 of @Rd16)→C			4						-	-	-	_	_	1	6
BOR #xx:3,@aa:8	В	C∨(#xx:3 of @aa:8)→C						4			-	-	_	_	_	‡	6
BIOR #xx:3,Rd	В	$C \lor (\overline{\#xx:3 \text{ of } Rd8}) \rightarrow C$		2							-	-	-	_	_	1	2

							essi icti											
Mnemonic	Operand size	Opera	tion	8/16		ı	@(d:16, Rn)	@-Rn/@Rn+	@aa: 8/16	@(d:8, PC)	aa	(Con	diti	on (cod	e	of states
	Ope			#xx:	Rn	@ Rn	p)@	@- F	@aa	p)@	@ @	Ι	Н	N	Z	V	C	No. 0
BIOR #xx:3,@Rd	В	Cv(#xx:3 of @Ro	d16) → C			4						-	_	_	_	-	‡	6
BIOR #xx:3,@aa:8	В	C∨(#xx:3 of @aa	:8)→C						4			-	-	-	_	-	‡	6
BXOR #xx:3,Rd	В	C⊕(#xx:3 of Rd8)→C		2							_	_	_	_	_	‡	2
BXOR #xx:3,@Rd	В	C⊕(#xx:3 of @Ro	d16)→C			4						-	-	_	-	_	1	6
BXOR #xx:3,@aa:8	В	C⊕(#xx:3 of @aa	::8) → C						4			_	_	_	_	_	1	6
BIXOR #xx:3,Rd	В	C⊕(#xx:3 of Rd8			2							_	_	_	_	_	1	2
BIXOR #xx:3,@Rd	В	C⊕(#xx:3 of @Re	d16) → C			4						_	_	_	_	_	1	6
BIXOR #xx:3,@aa:8	В	C⊕(#xx:3 of @aa	<u>i:8)</u> → C						4			_	-	_	-	_	1	6
BRA d:8 (BTd:8)	_	PC←PC+d:8								2		_	_	_	_	_	ı	4
BRN d:8 (BFd:8)	_	PC←PC+2								2		_	_	_	_	_	_	4
BHI d:8	-	if true then	C∨Z=0							2		_	_	_	_	-	_	4
BLS d:8	-	PC←PC+d:8	C∨Z=1							2		_	_	_	_	_	_	4
BCC d:8 (BHS d:8)	-	else next	C=0							2		-	_	-	-	-	_	4
BCS d:8 (BLO d:8)	-		C=1							2		-	-	-	-	-	_	4
BNE d:8	-		Z=0							2		_	-	_	-	-	-	4
BEQ d:8	-		Z=1							2		-	-	-	-	-	-	4
BVC d:8	-		V=0							2		-	-	-	-	-	_	4
BVS d:8	-		V=1							2		-	-	-	-	-	-	4
BPL d:8	-		N=0							2		_	_	-	_	-	_	4
BMI d:8	-		N=1							2		-	-	-	-	-	-	4
BGE d:8	-		N⊕V=0							2		_	-	_	-	_	-	4
BLT d:8	-		N⊕V=1							2		-	-	-	-	-	-	4
BGT d:8	-		$Z \lor (N \oplus V) = 0$							2		-	-	-	-	-	-	4
BLE d:8	-		Z∨(N⊕V)=1							2		_	_	_	-	-	-	4
JMP @Rn	-	PC←Rn16				2						-	-	-	-	-	-	4
JMP @aa:16	-	PC←aa:16							4			_	-	_	-	_	-	6
JMP @@aa:8	-	PC←@aa:8									2	-	-	-	-	-	-	8
BSR	-	SP-2→SP PC→ @SP PC← PC+d:8								2		-	-	-	_	-	-	6
JSR @Rn	-	SP-2→SP PC→@SP PC←Rn16				2						-	_	_	_	_	-	6

			i	Add inst	lres ruc	sin; tior	g m ı leı	ode ngtl	e/ 1								
rand size	Operation	8/16			:16, Rn)	kn/@Rn+	: 8/16	.8, PC)	aa	ied	(Con	diti	on	cod	e	No. of states
Ope		#xx	Rn	@Rr	p)@	@- R	@aa	(q)	@ @	Imp	I	Н	N	Z	V	C	No. 0
-	SP-2→SP						4				-	-	-	-	-	-	8
	$PC \rightarrow @SP$																
	PC←aa:16																
-	SP-2→SP								2		-	-	-	_	-	-	8
	PC→ @SP																
	PC←@aa:8																
-	PC←@SP									2	-	-	-	_	-	-	8
	SP+2→SP																
-	CCR←@SP									2	1	1	1	1	1	1	10
	SP+2→SP																
_	Transit to sleep mode.									2	_	_	-	_	_	_	2
В	#xx:8→CCR	2									1	1	‡	1	1	1	2
В	Rs8→CCR		2								1	1	1	1	1	1	2
В	CCR→ Rd8		2								_	_	_	_	_	_	2
В	CCR∧#xx:8→CCR	2									1	1	1	1	1	1	2
В	CCR∨#xx:8→CCR	2									1	1	1	1	1	1	2
В	CCR⊕#xx:8→CCR	2									1	1	1	‡	1	1	2
_	PC←PC+2									2	_	_	_	_	_	_	2
1	if R4L≠0 then Repeat @R5→@R6 R5+1→R5 R6+1→R6 R4L−1→R4L Until R4L=0									4	_	_	_	_	_	_	4
	B B B B B B -	- SP-2→SP PC→@SP PC→a:16 - SP-2→SP PC→@SP PC→@SP PC→@a:8 - PC←@SP SP+2→SP - CCR←@SP SP+2→SP PC←@SP SP+2→SP - Transit to sleep mode. B #xx:8→CCR B Rs8→CCR B CCR→Rd8 B CCR√#xx:8→CCR B CCR√#xx:8→CCR B CCR⊕#xx:8→CCR - PC←PC+2 - if R4L≠0 then Repeat @R5→@R6 R5+1→R5 R6+1→R6 R4L-1→R4L	- SP-2→SP PC→@SP PC→@SP PC→@SP PC→@SP PC→@SP PC→@SP PC→@SP SP+2→SP - CCR←@SP SP+2→SP - Transit to sleep mode. B #xx:8→CCR B CCR→Rd8 B CCR→Rd8 B CCR→Rd8 CCR→Rd8 B CCR→#xx:8→CCR CR→Rd8 B CCR→#xx:8→CCR CR→Rd8 CCR→Rd8 B CCR→#xx:8→CCR CR→Rd8 B CCR→#xx:8→CCR CR→Rd8 CR→Rd8 CR→Rd8 CR C	Operation SE Operation SE SE	Operation SE	SE Operation SE SE SE SE SE SE SE S	Operation Secondary Sec	Operation SP-2→SP	Operation Ope	- SP-2→SP PC→@SP PC→@SP PC→@aa:16 - SP-2→SP PC→@SP PC→@SP SP+2→SP - CCR←@SP SP+2→SP - Transit to sleep mode. B #xx:8→CCR B CCR→Rd8 CCR→Rd	SP-2→SP	Operation Operation	Operation Ope	Operation Ope	Operation Ope	SP-2-SP PC - @SP PC - @SP SP+2-SP PC - P	Operation Part Pa

Notes: The number of states is the number of states required for execution when the instructions and its operands are located in on-chip memory.

- ① Set to "1" when there is a carry or borrow from bit 11; otherwise cleared to "0."
- ② If the result is zero, the previous value of the flag is retained; otherwise the flag is cleared to "0."
- ③ Set to "1" if decimal adjustment produces a carry; otherwise cleared to "0."
- 4 The number of states is 4n+8 (n: initial value of R4L)
- ⑤ These instructions transfer data in synchronization with the E clock. The number of states varies depending on the synchronization delay.
- © Set to "1" if the divisor is negative; otherwise cleared to "0."
- ② Set to "1" if the divisor is zero: otherwise cleared to "0."

Appendix C. Number of Execution States

The tables in this appendix can be used to calculate the number of states required for instruction execution. Table C-1 indicates the number of states required for each cycle (instruction fetch, branch address read, stack operation, byte data access, word data access, internal operation). Table C-2 indicates the number of cycles of each type occurring in each instruction. The total number of states required for execution of an instruction can be calculated from these two tables as follows:

Execution states = $I \times S^I + J \times S^J + K \times S^K + L \times S^L + M \times S^M + N \times S^N$

Examples: Mode 1 (on-chip ROM disabled), stack located in external memory, 1 wait state inserted in external memory access.

1. BSET #0, @'FFC7

From table C-2:

$$I = L = 2$$
, $J = K = M = N = 0$

From table C-1:

$$S^{I} = 8$$
, $S^{L} = 3$

Number of states required for execution = $2 \times 8 + 2 \times 3 = 22$

2. JSR @@ 30

From table C-2:

$$I = 2$$
, $J = K = 1$, $L = M = N = 0$

From table C-1:

$$S^I \equiv S^J \equiv S^K \equiv 8$$

Number of states required for execution = $2 \times 8 + 1 \times 8 + 1 \times 8 = 32$

Table C-1. Number of States Taken by Each Cycle in Instruction Execution

Execution Status			Access Location	
(instruction cycle))	On-Chip Memory	On-Chip Reg. Field	External Memory
Instruction fetch	SI			
Branch address read	S_1		6	6 + 2m
Stack operation	SK	2		
Byte data access	SL		3	3 + m
Word data access	SM		6	6 + 2m
Internal operation	SN		2	

Notes:

- 1. m: Number of wait states inserted in access to external device.
- 2. The byte data access cycle to an external device by the MOVFPE and MOVTPE instructions requires 9 to 16 states since it is synchronized with the E clock. See the *Hardware Manual* for timing details.

Table C-2. Number of Cycles in Each Instruction

Instruction	Mnemonic	Instruction Fetch	Branch Addr. Read	Stack Operation	Byte Data Access	Word Data Access	Internal Operation
		I	J	K	L	М	N
ADD	ADD.B #xx:8, Rd	1					
	ADD.B Rs, Rd	1					
	ADD.W Rs, Rd	1					
ADDS	ADDS.W #1/2, Rd	1					
ADDX	ADDX.B #xx:8, Rd	1					
	ADDX.B Rs, Rd	1					
AND	AND.B #xx:8, Rd	1					
	AND.B Rs, Rd	1					
ANDC	ANDC #xx:8, CCR	1					
BAND	BAND #xx:3, Rd	1					
	BAND #xx:3, @Rd	2			1		
	BAND #xx:3, @aa:8	2			1		
Всс	BRA d:8 (BT d:8)	2					
	BRN d:8 (BF d:8)	2					
	BHI d:8	2					
	BLS d:8	2					
	BCC d:8 (BHS d:8)	2					
	BCS d:8 (BLO d:8)	2					
	BNE d:8	2					
	BEQ d:8	2					
	BVC d:8	2					
	BVS d:8	2					
	BPL d:8	2					
	BMI d:8	2					
	BGE d:8	2					
	BLT d:8	2					
	BGT d:8	2					
	BLE d:8	2					
BCLR	BCLR #xx:3, Rd	1					
	BCLR #xx:3, @Rd	2			2		
	BCLR #xx:3, @aa:8	2			2		
	BCLR Rn, Rd	1					

Instruction	Mnemonic	Instruction Fetch	Branch Addr. Read		Byte Data Access		Internal Operation
		I	J	K	L	M	N
BCLR	BCLR Rn, @Rd	2			2		
	BCLR Rn, @aa:8	2			2		
BIAND	BIAND #xx:3, Rd	1					
	BIAND #xx:3, @Rd	2			1		
	BIAND #xx:3, @aa:8	2			1		
BILD	BILD #xx:3, Rd	1					
	BILD #xx:3, @Rd	2			1		
	BILD #xx:3, @aa:8	2			1		
BIOR	BIOR #xx:3, Rd	1					
	BIOR #xx:3, @Rd	2			1		
	BIOR #xx:3, @aa:8	2			1		
BIST	BIST #xx:3, Rd	1					
	BIST #xx:3, @Rd	2			2		
	BIST #xx:3, @aa:8	2			2		
BIXOR	BIXOR #xx:3, Rd	1					
	BIXOR #xx:3, @Rd	2			1		
	BIXOR #xx:3, @aa:8	2			1		
BLD	BLD #xx:3, Rd	1					
	BLD #xx:3, @Rd	2			1		
	BLD #xx:3, @aa:8	2			1		
BNOT	BNOT #xx:3, Rd	1					
	BNOT #xx:3, @Rd	2			2		
	BNOT #xx:3, @aa:8	2			2		
	BNOT Rn, Rd	1					
	BNOT Rn, @Rd	2			2		
	BNOT Rn, @aa:8	2			2		
BOR	BOR #xx:3, Rd	1					
	BOR #xx:3, @Rd	2			1		
	BOR #xx:3, @aa:8	2			1		
BSET	BSET #xx:3, Rd	1					
	BSET #xx:3, @Rd	2			2		
	BSET #xx:3, @aa:8	2			2		
	BSET Rn, Rd	1					
	BSET Rn, @Rd	2			2		

Instruction	Mnemonic	Instruction Fetch	Branch Addr. Read	Stack Operation	Byte Data Access	Word Data Access	Internal Operation
		I	J	K	L	M	N
BSET	BSET Rn, @aa:8	2			2		
BSR	BSR d:8	2		1			
BST	BST #xx:3, Rd	1					
	BST #xx:3, @Rd	2			2		
	BST #xx:3, @aa:8	2			2		
BTST	BTST #xx:3, Rd	1					
	BTST #xx:3, @Rd	2			1		
	BTST #xx:3, @aa:8	2			1		
	BTST Rn, Rd	1					
	BTST Rn, @Rd	2			1		
	BTST Rn, @aa:8	2			1		
BXOR	BXOR #xx:3, Rd	1					
	BXOR #xx:3, @Rd	2			1		
	BXOR #xx:3, @aa:8	2			1		
CMP	CMP. B #xx:8, Rd	1					
	CMP. B Rs, Rd	1					
	CMP.W Rs, Rd	1					
DAA	DAA.B Rd	1					
DAS	DAS.B Rd	1					
DEC	DEC.B Rd	1					
DIVXU	DIVXU.B Rs, Rd	1					6
EEPMOV	EEPMOV	2			2n+2*1		
INC	INC.B Rd	1					
JMP	JMP @Rn	2					
	JMP @aa:16	2					1
	JMP @@aa:8	2	1				1
JSR	JSR @Rn	2		1			
	JSR @aa:16	2		1			1
	JSR @@aa:8	2	1	1			
LDC	LDC #xx:8, CCR	1					
	LDC Rs, CCR	1					
MOV	MOV.B #xx:8, Rd	1					
	MOV.B Rs, Rd	1					
	MOV.B @Rs, Rd	1			1		

Instruction	Mnemonic	Instruction Fetch	Branch Addr. Read	Stack Operation	Byte Data Access	Word Data Access	Internal Operation
		I	J	K	L	M	N
MOV	MOV.B @(d:16, Rs), Rd	2			1		
	MOV.B @Rs+, Rd	1			1		1
	MOV.B @aa:8, Rd	1			1		
	MOV.B @aa:16, Rd	2			1		
	MOV.B Rs, @Rd	1			1		
	MOV.B Rs, @(d:16, Rd)	2			1		
	MOV.B Rs, @-Rd	1			1		1
	MOV.B Rs, @aa:8	1			1		
	MOV.B Rs, @aa:16	2			1		
	MOV.W #xx:16, Rd	2					
	MOV.W Rs, Rd	1					
	MOV.W @Rs, Rd	1				1	
	MOV.W @(d:16, Rs), Rd	2				1	
	MOV.W @Rs+, Rd	1				1	1
	MOV.W @aa:16, Rd	2				1	
	MOV.W Rs, @Rd	1				1	
	MOV.W Rs, @(d:16, Rd)	2				1	
	MOV.W Rs, @-Rd	1				1	1
	MOV.W Rs, @aa:16	2				1	
MOVFPE	MOVFPE @aa:16, Rd	2			1*2		
MOVTPE	MOVTPE Rs, @aa:16	2			1*2		
MULXU	MULXU.B Rs, Rd	1					6
NEG	NEG.B Rd	1					
NOP	NOP	1					
NOT	NOT.B Rd	1					
OR	OR.B #xx:8, Rd	1					
	OR.B Rs, Rd	1					
ORC	ORC #xx:8, CCR	1					
ROTL	ROTL.B Rd	1					
ROTR	ROTR.B Rd	1					
ROTXL	ROTXL.B Rd	1					
ROTXR	ROTXR.B Rd	1					
RTE	RTE	2		2			1
RTS	RTS	2		1			1

Instruction	Mnemonic	Instruction Fetch	Branch Addr. Read		Byte Data Access	Word Data Access	Internal Operation
		I	J	K	L	M	N
SHAL	SHAL.B Rd	1					
SHAR	SHAR.B Rd	1					
SHLL	SHLL.B Rd	1					
SHLR	SHLR.B Rd	1					
SLEEP	SLEEP	1					
STC	STC CCR, Rd	1					
SUB	SUB.B Rs, Rd	1					
	SUB.W Rs, Rd	1					
SUBS	SUBS.W #1/2, Rd	1					
SUBX	SUBX.B #xx:8, Rd	1					
	SUBX.B Rs, Rd	1					
XOR	XOR.B #xx:8, Rd	1					
	XOR.B Rs, Rd	1					
XORC	XORC #xx:8, CCR	1					

Notes:

- *1 n: Initial value in R4L. The source and destination operands are accessed n+1 times each.
- *2 Data access requires 9 to 16 states.

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